

# Logic Gates

CLASS 9

## HW 6.2-solution (Continued):

10) DeMorgan's law:

$$(x + y)' = x' y'$$

$$(x y)' = x' + y' \quad \forall x, y \in B$$

### Proof

By duality we need to prove only the first equality:  $(x + y)' = x' y'$

We use the same reasoning as for the previous property, 9).

That is, we use the axiom that defines the complement, namely 5):

In the first equality, which is the element and which do we have to prove is the complement?

Element =  $x + y$ , and its complement =  $x' y'$ .

Substitute in 5). We need to prove:

$$x + y + x' y' = 1$$

$$(x + y) \cdot (x' y') = 0$$

HW 7.1 - assigned: If not done yet, continue by proving that the two equalities above hold!

$$x + y + x' y' = 1$$

$$(x + y) \cdot (x' y') = 0$$

$$(x + y) \cdot 1 + x' y' =$$

$$= (x + y) \cdot (x + x') + x' y' =$$

$$= xx + xx' + yx + yx' + x' y' =$$

$$= (x + yx) + (xx') + (yx' + x' y') =$$

$$= x + 0 + x' (y + y') =$$

$$= x + x' \cdot 1 = x + x' = 1$$

2 Identity

7 Idempotent

12 Absorbtion, 5 Complement,

5 Complement, 2 Identity

$$(x + y) \cdot (x' y') =$$

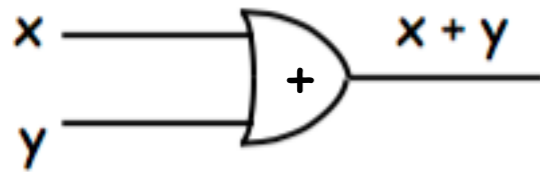
$$= xx' y + yx' y' =$$

5 Complement

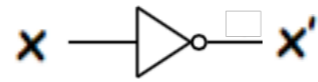
$$= 0 \cdot y + 0 \cdot x' = 0$$

8 Identity, Annulment

# Gates



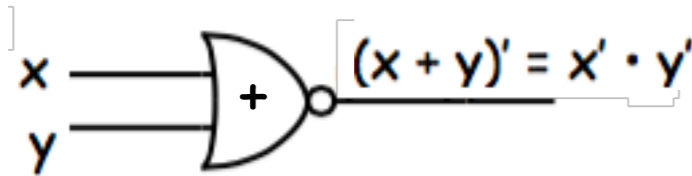
**OR**



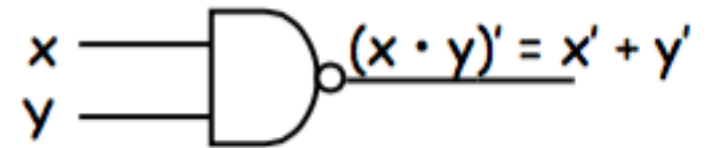
**Inverter**



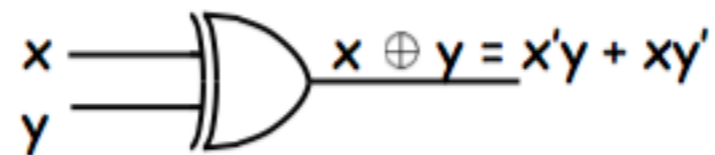
**AND**



**NOR**



**NAND**



**XOR**  
**Exclusive OR**

## Example - logic diagram

Remember:

AND:

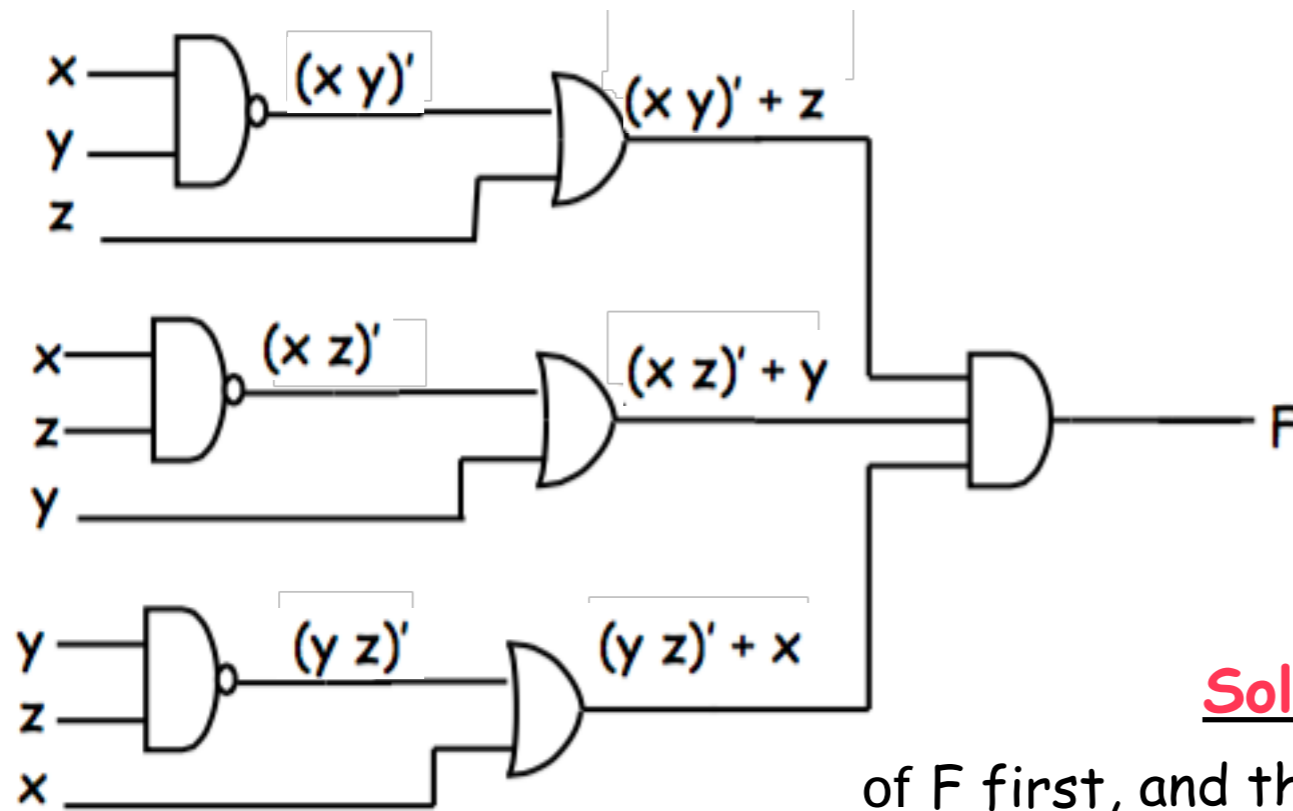
|   |   |   |
|---|---|---|
| • | 0 | 1 |
| 0 | 0 | 0 |
| 1 | 0 | 1 |

OR:

|   |   |   |
|---|---|---|
| + | 0 | 1 |
| 0 | 0 | 1 |
| 1 | 1 | 1 |

Problem:

Find the truth table and expression of the function F.



Solution

I recommend finding the expression of F first, and then deduce the truth table from the expression.

We have:  $(xy)' + z = x' + y' + z$

$(xz)' + y = x' + z' + y$

$(yz)' + x = y' + z' + x$

What is then F?

$$F = (x' + y' + z)(x' + z' + y)(y' + z' + x)$$

What can we say about this product?

It is mostly = 0, like every product.

It is = 0 when any one of its factors = 0

$$x' + y' + z = 0 \iff \begin{matrix} x'=0 \text{ and} \\ y'=0 \text{ and} \\ x=1 \text{ and} \\ y=1 \text{ and} \\ z=0 \end{matrix}$$

Similarly we have:

$$x' + y + z' = 0 \iff \begin{matrix} x=1 \text{ and} \\ y=0 \text{ and} \\ z=1 \end{matrix}$$

$$x + y' + z' = 0 \iff \begin{matrix} x=0 \text{ and} \\ y=1 \text{ and} \\ z=1 \end{matrix}$$

In rest, for all other values of x, y, z, F = 1

| x | y | z | F |
|---|---|---|---|
| 0 | 0 | 0 | 1 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 0 |
| 1 | 1 | 1 | 1 |

Truth table

## HW 8.1

### Partial Solution + Hint

#### (a) Prove: Inhibition is not commutative

Table gives two functions:

|             |       |            |                   |
|-------------|-------|------------|-------------------|
| $F_2 = xy'$ | $x/y$ | Inhibition | $x$ , but not $y$ |
| $F_4 = x'y$ | $y/x$ | Inhibition | $y$ , but not $x$ |

Let's choose  $F_2 = x/y = x y'$

$$\begin{array}{l} \text{F}_2 \text{ commutative} \stackrel{\text{Det}}{\Leftrightarrow} \boxed{F_2(x, y) = F_2(y, x) \quad \forall x, y} \stackrel{\text{Det}}{\Leftrightarrow} \boxed{x/y = y/x \quad \forall x, y} \\ \text{F}_2 \text{ not commutative} \stackrel{\text{Det}}{\Leftrightarrow} \exists x, y: F_2(x, y) \neq F_2(y, x) \stackrel{\text{Det}}{\Leftrightarrow} \exists x, y: x/y \neq y/x \end{array}$$

In other words we need to find a counterexample to  $x/y = y/x$ , that is find values in  $\{0, 1\}$  for  $x, y$ , such that we have  $x/y \neq y/x \stackrel{\text{Det}}{\Leftrightarrow} xy' \neq yx'$  as every B.A. has 0 and 1 as elements!

Make one side 1 and at the same time the other side 0.

$x = 1$  and  $y = 0$  imply:

$$LS = 1 \cdot 1 = 1 \quad RS = 0 \cdot 0 = 0 \quad \rightarrow LS \neq RS \quad \rightarrow \text{F}_2 \text{ is not commutative}$$

Done

HW 8.1 - continued Partial Solution + Hint - continued

(b) Prove: Exclusive OR (XOR) is associative Table:

$$F_6 = xy' + x'y \quad x \oplus y \quad \text{Exclusive-OR} \quad x \text{ or } y, \text{ but not both}$$

$$F_6(x, y) = xy' + x'y = x \oplus y$$

$$F_6 \text{ associative} \stackrel{\text{Det}}{\Leftrightarrow} (x \oplus y) \oplus z = x \oplus (y \oplus z) \quad \forall x, y, z$$

$$\stackrel{\text{Det}}{\Leftrightarrow} F_6(F_6(x, y), z) = F_6(x, F_6(y, z)) \quad \forall x, y, z$$

We need to prove this for all BAs, so we need to use the axioms and properties of a BA.

$$LS = (xy' + x'y) \oplus z = (xy' + x'y)z' + (xy' + x'y)'z =$$

$$= xy'z' + x'yz' + (x' + y)(x + y')z = xy'z' + x'yz' + x'y'z + xyz \quad \text{5 Complement}$$

$$RS = x \oplus (yz' + y'z) = x(yz' + y'z)' + x'(yz' + y'z) = x(y' + z)(y + z') + x'yz' + x'y'z =$$

$$= xy'z' + xyz + x'yz' + x'y'z \quad \rightarrow LS = RS \quad \rightarrow \underline{F_6 \text{ is associative}}$$

Done

### HW 7.3 (continuation from class 8)

2-9 b)

$$\overline{((\bar{v}w + x)y + \bar{z})} = \overline{((v + \bar{w}) \cdot \bar{x} + \bar{y})z} =$$

$$= (v\bar{x} + \bar{w}\bar{x} + \bar{y})z = \underline{v\bar{x}z + \bar{w}\bar{x}z + \bar{y}z}$$

c)  $\overline{(w \times (\bar{y}z + y\bar{z}) + \bar{w}\bar{x}(\bar{y} + z)(y + \bar{z}))} =$

$$= (\bar{w} + \bar{x} + (y + \bar{z})(\bar{y} + z))(w + x + y\bar{z} + \bar{y}z) =$$

$$= (\bar{w} + \bar{x} + yz + \bar{y}\bar{z})(w + x + y\bar{z} + \bar{y}z) =$$

$$= \bar{w}x + \bar{w}y\bar{z} + \bar{w}\bar{y}z + \bar{x}w + \bar{x}y\bar{z} + \bar{x}\bar{y}z + wyz + xyz +$$
$$+ w\bar{y}\bar{z} + x\bar{y}\bar{z}$$



## HW 7.3 (continuation)

2-9 d)

$$\begin{aligned} & \overbrace{\left( (A + \bar{B} + c)(\bar{A}\bar{B} + c)(A + \bar{B}\bar{c}) \right)} = \\ & = \bar{A}B\bar{c} + (A+B)\cdot\bar{c} + \bar{A}\cdot(B+c) \quad (4) \\ & = \bar{A}B\bar{c} + A\bar{c} + B\bar{c} + \bar{A}B + \bar{A}c \quad (12) \\ & = A\bar{c} + B\bar{c} + \bar{A}B + \bar{A}c \end{aligned}$$

What now? Expand and then contract again:

$$\begin{aligned} & = A\bar{B}\bar{c} + A\bar{B}\bar{c} + \cancel{AB\bar{c}} + \cancel{\bar{A}B\bar{c}} + \bar{A}B\bar{c} + \cancel{\bar{A}B\bar{c}} + \\ & + \cancel{\bar{A}Bc} + \bar{A}\bar{B}c = \underline{A\bar{B}\bar{c}} + \underline{AB\bar{c}} + \underline{\bar{A}B\bar{c}} + \\ & + \underline{\bar{A}Bc} + \underline{\bar{A}\bar{B}c} = \underline{A\bar{c}} + \underline{\bar{A}B} + \underline{\bar{A}c} \end{aligned}$$