

MIDTERM REVIEW

Review

Topics

(1) Number Representations

- $(r - 1)$'s, r 's complements, algorithm for subtraction, etc
- Find basis x if $(\dots)_x = \dots$

(2) (3) - Boolean algebra:

- Expression manipulation
- Boolean operators

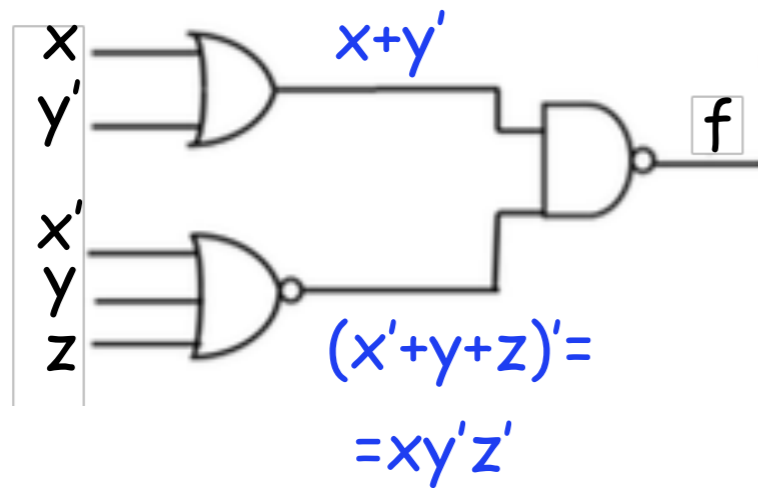
(4) distribution (both)
(12) absorption
(10) De Morgan

(4) - Circuit - logic diagram (gates)

B. functions; truth table

(4) Gates

Find the truth table + expression of f



$$\begin{aligned} f &= ((x+y')xy'z')' = \\ &= (x+y')' + x' + y + z = \\ &= x'y + x' + y + z = \boxed{x' + y + z} \end{aligned}$$

$x' (12)$

x	y	z	f
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

What can we say about the truth table for f?

When do we have $f = 0$, and when is $f = 1$?

It is 0 exactly when $x'=0$, and $y=0$, and $z=0$, that is for the triple $x y z = 1 0 0$.

In rest we have $f = 1$.

(3) Operators

HW 8.1

(a) Prove: Inhibition is not associative

We chose $F_2 = x/y = x y'$

$$F_2 \text{ associative} \stackrel{\text{Det}}{\Leftrightarrow} F_2 (F_2 (x, y), z) = F_2 (x, F_2 (y, z)) \quad \forall x, y, z$$



$$F_2 (xy', z) = F_2 (x, yz')$$



$$xy'z' = x(yz)'$$



$$xy'z' = x(y'+z) \quad \forall x, y, z \quad (10)$$

$$F_2 \text{ not associative} \stackrel{\text{Det}}{\Leftrightarrow} \exists x, y, z: xy'z' \neq x(y'+z)$$

Make, say, $RS = 1$, while $LS = 0$: We need: $x = 1$ and ($y' = 1$ or $z = 1$)

There are multiple such counterexamples.

We only need one, for instance $x = 1, z = 1$, and any value for y , say $y = 0$.

We have $LS = 0 \neq 1 = RS$.

We just proved: **F_2 is not associative**

(3) Operators - continued

HW 8.1 - continued

(b) Prove: Exclusive OR (XOR) is commutative

$$F_6(x, y) = x y' + x' y = x \oplus y$$

$$F_6 \text{ commutative} \stackrel{\text{Det}}{\iff} F_6(x, y) = F_6(y, x) \quad \forall x, y$$

$$\oplus \text{ commutative} \stackrel{\text{Det}}{\iff} x \oplus y = y \oplus x \quad \forall x, y$$

$$\underline{x y'} + \underline{x' y} = \underline{y x'} + \underline{y' x} \quad \forall x, y$$

The equality holds for both binary operations + and · are commutative.

We just proved: F_6 , that is Exclusive OR (XOR) is commutative

(2) Expression manipulation

Basically what we have been doing all along - old HW

2.6)

$$\text{a) } \overbrace{(xy' + x'y)'}^{(10)} = \underbrace{(x' + y) \cdot (x + y')} = \boxed{xy + x'y'}$$

$$\text{b) } \overbrace{((AB' + C)D' + E)'}^{(10)} = \underbrace{((A' + B) \cdot C' + D)} \cdot E' = \boxed{(A' + B + D)(C' + D)E'}^{(4)}$$

$$\text{c) } \overbrace{((x + y' + z)(x' + z')(x + y))'}^{(10)} = \boxed{x' \cdot y \cdot z' + x \cdot z + x' \cdot y'}$$

More exercises:

$$\text{A. } \underbrace{xyz' + x'yz}_{(4)} + \underbrace{xyz + x'yz'}_{(5)} = \underbrace{xy}_{(4)} + \underbrace{yz}_{(5)} + \underbrace{x'y}_{(2)} = \underbrace{y}_{(5)} + \underbrace{yz}_{(12)} = \boxed{y}$$

$xy(z+z') = xy$
 $(x+x')yz$

$$\begin{aligned}
 \text{B. } & \underbrace{xyzt + x'yt + xyz't + yzt}_{xyt} \stackrel{(4)}{=} \underbrace{xyt + x'yt}_{yt} + yzt \stackrel{(4)}{=} yt + yzt \stackrel{(12)}{=} \boxed{yt} \\
 & \quad \quad \quad (5) \quad \quad \quad (5) \quad \quad \quad (2)
 \end{aligned}$$

C. Prove: $xy + x'z + yz = xy + x'z$

Expand the terms that are in $LS \setminus RS$, namely yz

$$\begin{aligned}
 \boxed{LS} &= xy + x'z + (x + x')yz \stackrel{(12) \times 2}{=} \underbrace{xy}_{\text{red}} + \underbrace{x'z}_{\text{blue}} + \underbrace{xyz}_{\text{red}} + \underbrace{x'yz}_{\text{blue}} = \\
 &= \underbrace{xy}_{\text{red}} + \underbrace{x'z}_{\text{blue}} = \boxed{RS}
 \end{aligned}$$

(1) Number Representations

Review conversions and operations in any base.

Let's do the r 's / $(r-1)$'s complements and apply the algorithm for subtracting two numbers.

1. $r = 16$

Compute the $(r-1)$'s complement of $3ABF09_{\text{hex}}$. We have $r - 1 = 16 - 1 = 15$. Subtract each digit from $F=15$:

$$(3ABF09_{\text{hex}})_{15'C} = \boxed{C540F6}$$

Let's do it via $r = 2$ (all computations are in binary). How?

$$3ABF09_{\text{hex}} = 0011\ 1010\ 1011\ 1111\ 0000\ 1001_2 \quad \begin{array}{l} \text{1's complement} \\ \text{-----}\rightarrow \end{array}$$

$$\begin{array}{l} \text{1's complement} \\ \text{-----}\rightarrow \end{array} 1100\ 0101\ 0100\ 0000\ 1111\ 0110_2 = \boxed{C540F6_{\text{hex}}}$$

Decimal \longleftrightarrow (2421) code

(8, 4, -2, -1) code

Yes, for self-complementary

Remember:

Decimal $\not\longleftrightarrow$ BCD

No, not self-complementary

Table

Four Different Binary Codes for the Decimal Digits

Decimal Digit	BCD 8421	2421	Excess-3	8, 4, -2, -1
0	0000	0000	0011	0000
1	0001	0001	0100	0111
2	0010	0010	0101	0110
3	0011	0011	0110	0101
4	0100	0100	0111	0100
5	0101	1011	1000	1011
6	0110	1100	1001	1010
7	0111	1101	1010	1001
8	1000	1110	1011	1000
9	1001	1111	1100	1111
Unused bit combinations	1010	0101	0000	0001
	1011	0110	0001	0010
	1100	0111	0010	0011
	1101	1000	1101	1100
	1110	1001	1110	1101
	1111	1010	1111	1110

