

# Stack Implementations

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# Today's Plan



Stack Implementations:

Array

Vector

Linked Chain

# Stack ADT

```
#ifndef STACK_H_
#define STACK_H_

template<typename ItemType>
class Stack
{
public:
    Stack();
    void push(const ItemType& new_entry); //adds an element to top of stack
    void pop(); // removes element from top of stack
    ItemType top() const; // returns a copy of element at top of stack
    int size() const; // returns the number of elements in the stack
    bool isEmpty() const; //returns true if no elements on stack, else false

private:
    //implementation details here

}; //end Stack

#include "Stack.cpp"
#endif // STACK_H_`
```

# ADT vs Data Structure

**ADT** is the logical/abstract description of the organization and operations on the data

**Data Structure** is the representation/implementation of the ADT

We may have multiple implementations of the same ADT

- 1 ADT
- Multiple Data Structures

To complicate matters, a data structure may be implemented using other data structures

- stack implemented using vector
- priority queue implemented using heap (more on this later)

If main() instantiates a stack, it is using a stack data structure, no matter which implementation we choose

# Choose a Data Structure

Array?

Vector?

Linked chain?

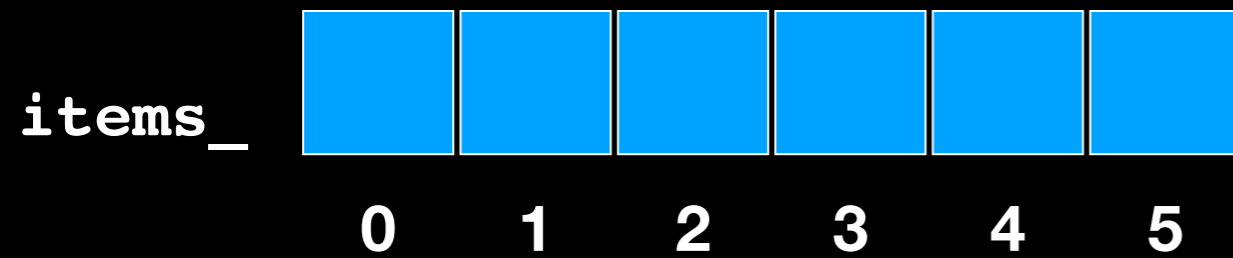
# Choose a Data Structure

Inserting and removing from same end (**LIFO**)

Goal: minimize work - Ideally  $O(1)$

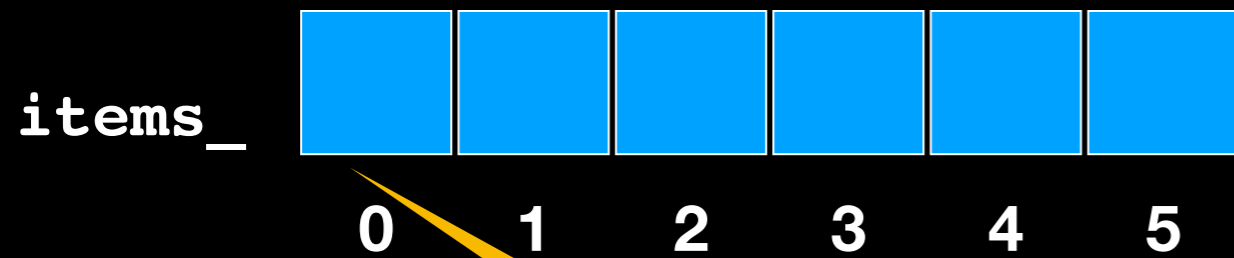
What would you suggest?

# Array



Where is the top  
of the stack?

# Array



```
item_count_ = 0
```

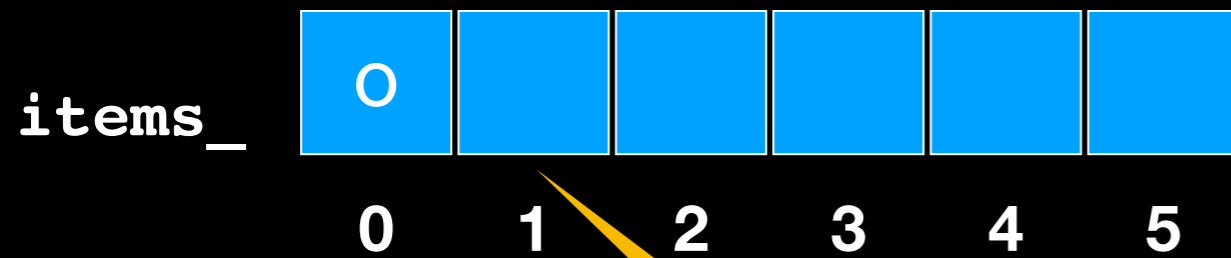
```
max_items_ = 6
```

Top of the stack:  
`items_[item_count_]`



# Array

```
push('0')
```



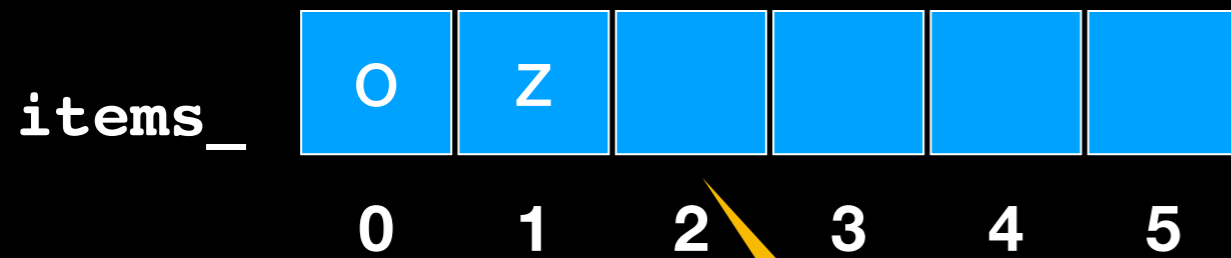
```
item_count_ = 1
```

```
max_items_ = 6
```

Top of the stack:  
`items_[item_count_]`

# Array

```
push('Z')
```



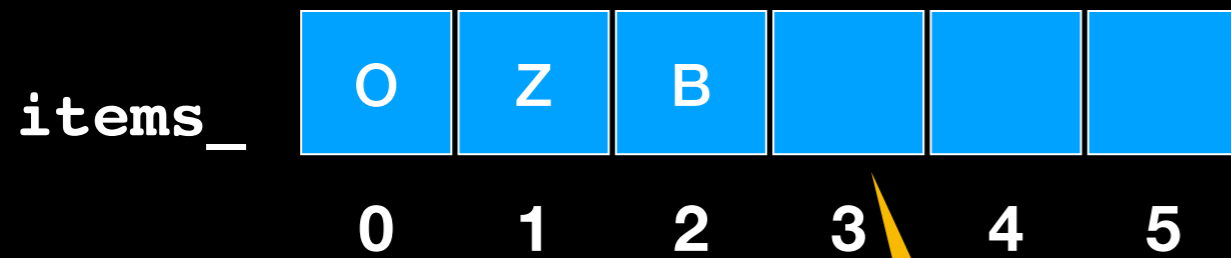
```
item_count_ = 2
```

```
max_items_ = 6
```

Top of the stack:  
`items_[item_count_]`

# Array

`push('B')`

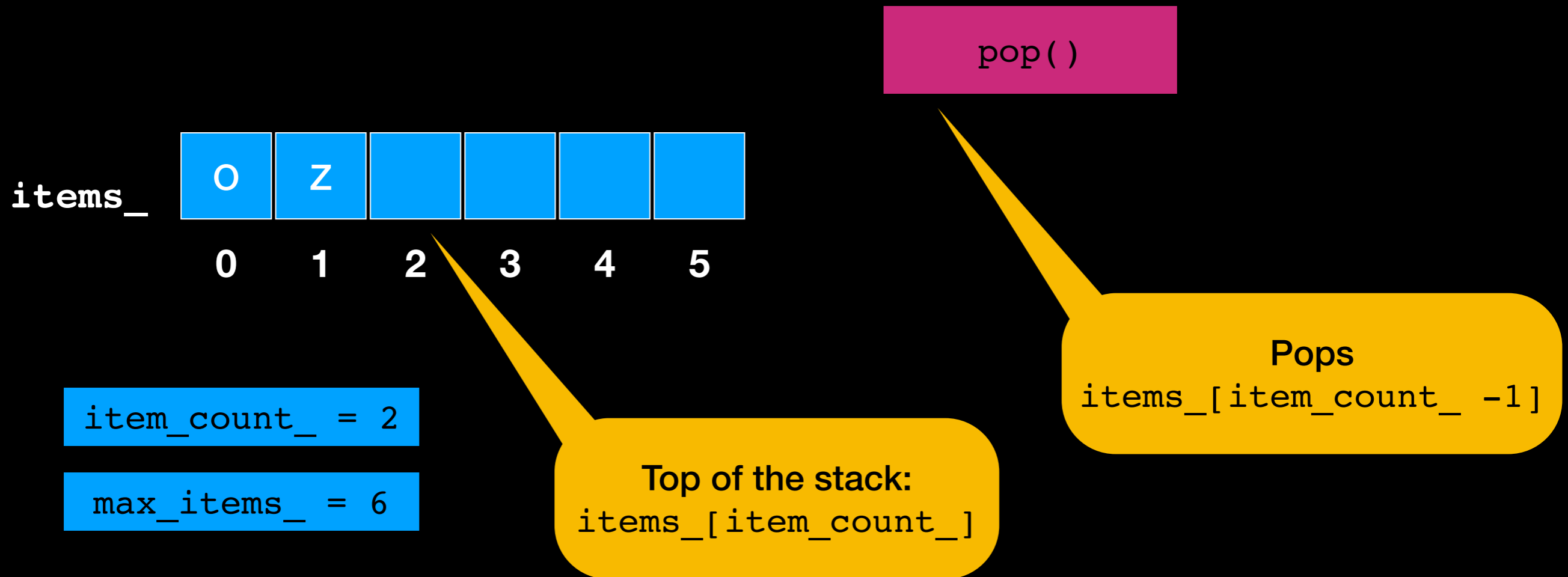


`item_count_ = 3`

`max_items_ = 6`

Top of the stack:  
`items_[item_count_]`

# Array



# Array Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

*push*:  $O(1)$

*pop* :  $O(1)$

*top* :  $O(1)$

GREAT!!!!

# Array Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

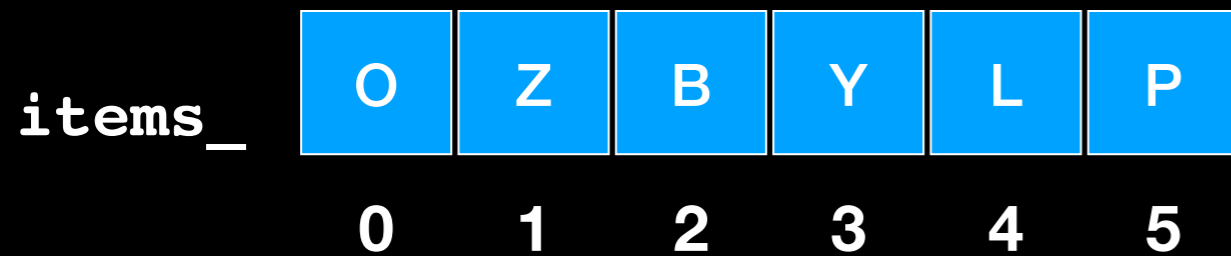
*push*:  $O(1)$

*pop* :  $O(1)$

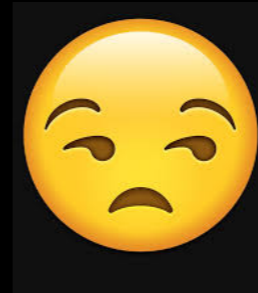
*top* :  $O(1)$

GREAT???

# Array



push('T')



Sorry Stack is Full!!!

item\_count\_ = 6

max\_items\_ = 6

Top of the stack:  
items\_[item\_count\_]

# Vector

```
std::vector<T> some_vector;
```

So what is a vector really?



# Vector

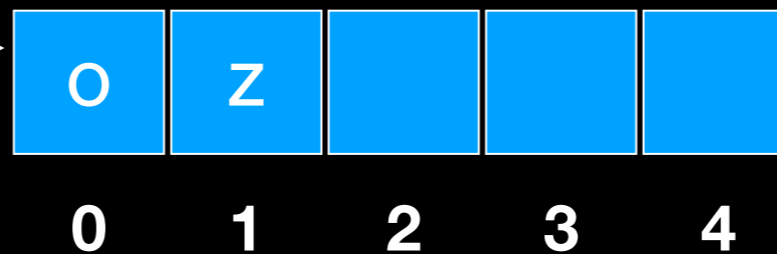
```
std::vector<T> some_vector;
```

So what is a vector really?

Push and pop same as  
with arrays

**Vector (simplified)**

buffer\_ =  
len\_ = 2  
capacity\_ = 5



# Vector

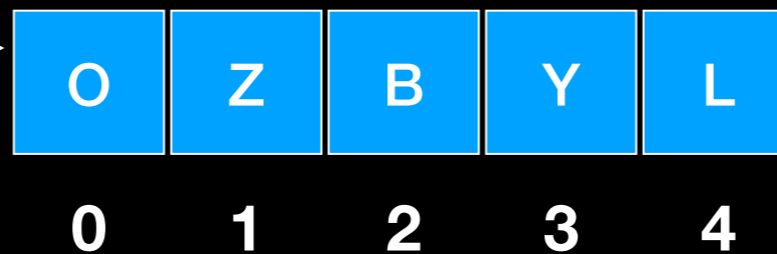
```
std::vector<T> some_vector;
```

So what is a vector really?

Stack is Full?

**Vector (simplified)**

buffer\_ =  
len\_ = 5  
capacity\_ = 5



# Vector

```
std::vector<T> some_vector;
```

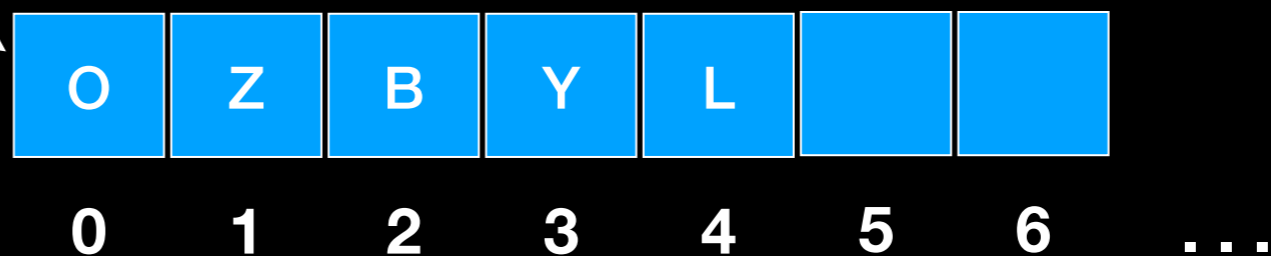
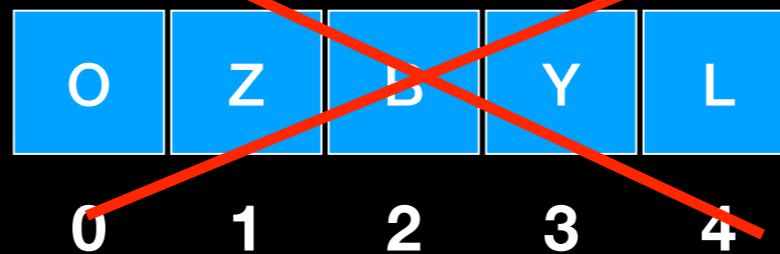
So what is a vector really?



No, I'll Grow!!!

**Vector (simplified)**

buffer\_ =  
len\_ = 5  
capacity\_ = ?



# Lecture Activity

How much should it grow?

Write a short paragraph arguing the **pros** and **cons** of growing by the amount you propose

# Vector Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

*push*:  $O(1)$

*pop* :  $O(1)$

*top* :  $O(1)$

GREAT!!!!

# Vector Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

*push*:  $O(1)$

*pop* :  $O(1)$

*top* :  $O(1)$

GREAT???

Except when stack is full must:

- allocate new array
- copy elements in new array
- delete old array

# Vector Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

*push*:  $O(1)$

*pop* :  $O(1)$

*top* :  $O(1)$

GREAT???

Except when stack is full must:

- allocate new array  $O(1)$
- copy elements in new array  $O(n)$
- delete old array  $O(1)$

# How should Vector grow?

Sometimes 1 "step"

Sometimes  $n$  "steps"

Consider behavior **over several pushes**



# Vector Growth: a naive approach

```
std::vector<T> some_vector;
```

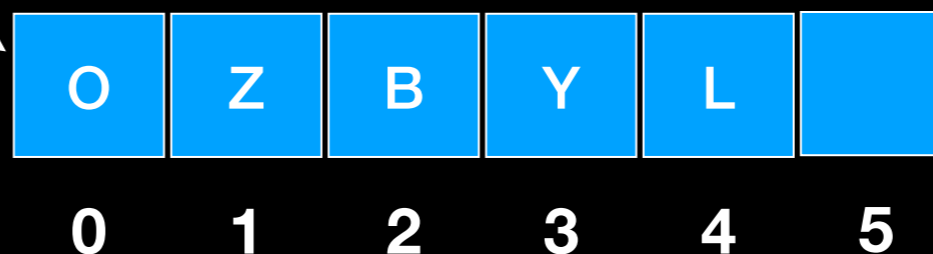
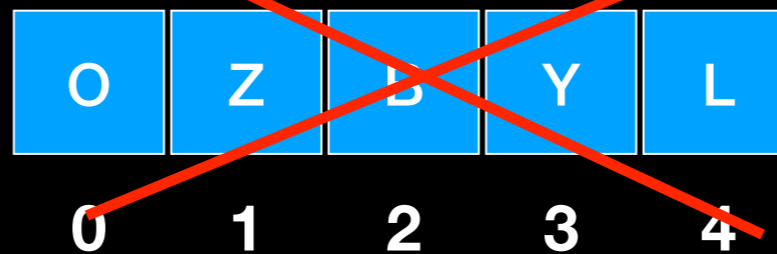
So what is a vector really?



I'll Grow!!!  
I will add space for the  
item to be added

**Vector (simplified)**

buffer\_ =  
len\_ = 5  
capacity\_ = 6



# Vector Growth: a naive approach

If vector **grows by 1 each time**, every push costs  $n$  "steps"

Cost of pushes:

$$1 + 2 + 3 + 4 + 5 + \dots + n$$
$$= n(n+1)/2$$

# Vector Growth: a naive approach

If vector **grows by 1 each time**, every push costs  $n$  "steps"

**Cost of  $n$  pushes:**

$$1 + 2 + 3 + 4 + 5 + \dots + n$$

$$= n(n+1)/2$$

$$= n^2/2 + n/2 \quad O(n^2)$$

# Vector Growth: a better approach

```
std::vector<T> some_vector;
```

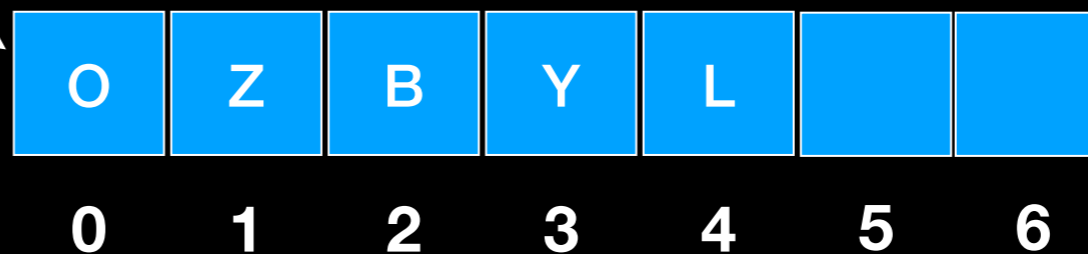
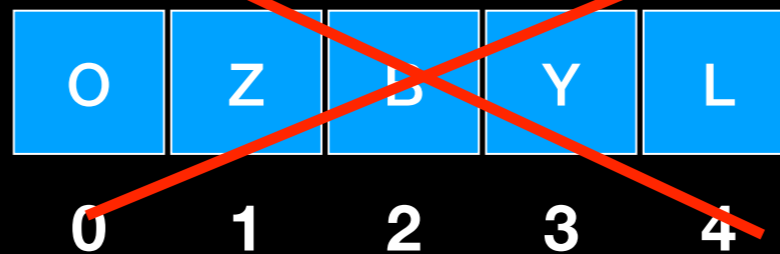
So what is a vector really?



I'll Grow!!!  
I will add two more slots!

**Vector (simplified)**

buffer\_ =  
len\_ = 5  
capacity\_ = 7



# Vector Growth: a better approach

If vector **grows by 2 each time**,

Let a “**hard push**” be one where the whole vector needs to be copied

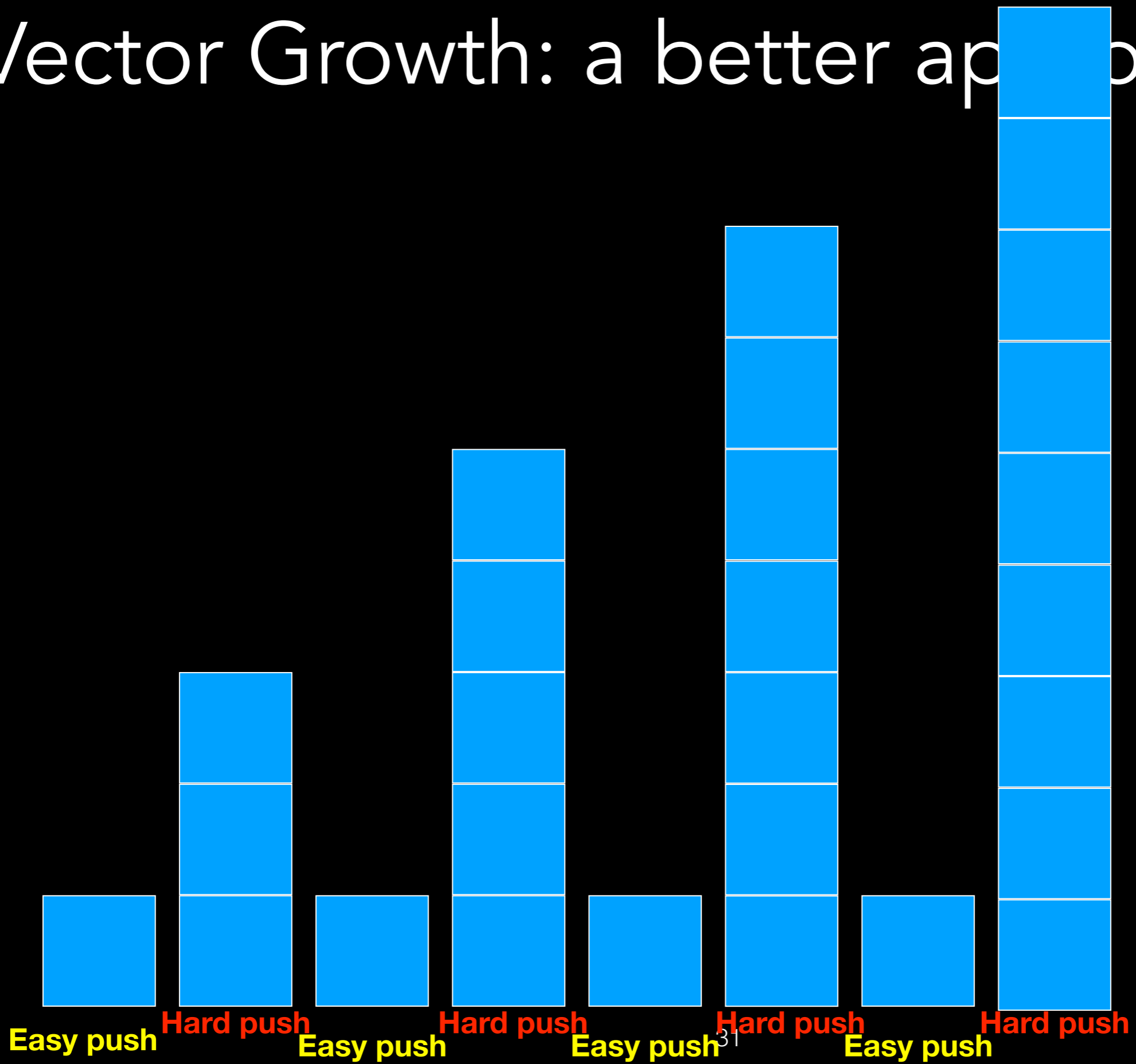
When vector is not copied we have an “**easy push**”

Now **half** our pushes will be **easy (1 step)** and **half** will be **hard (n steps)**

So if reconsider the work **over several pushes?**

Analysis visualization adapted from Keith Schwarz

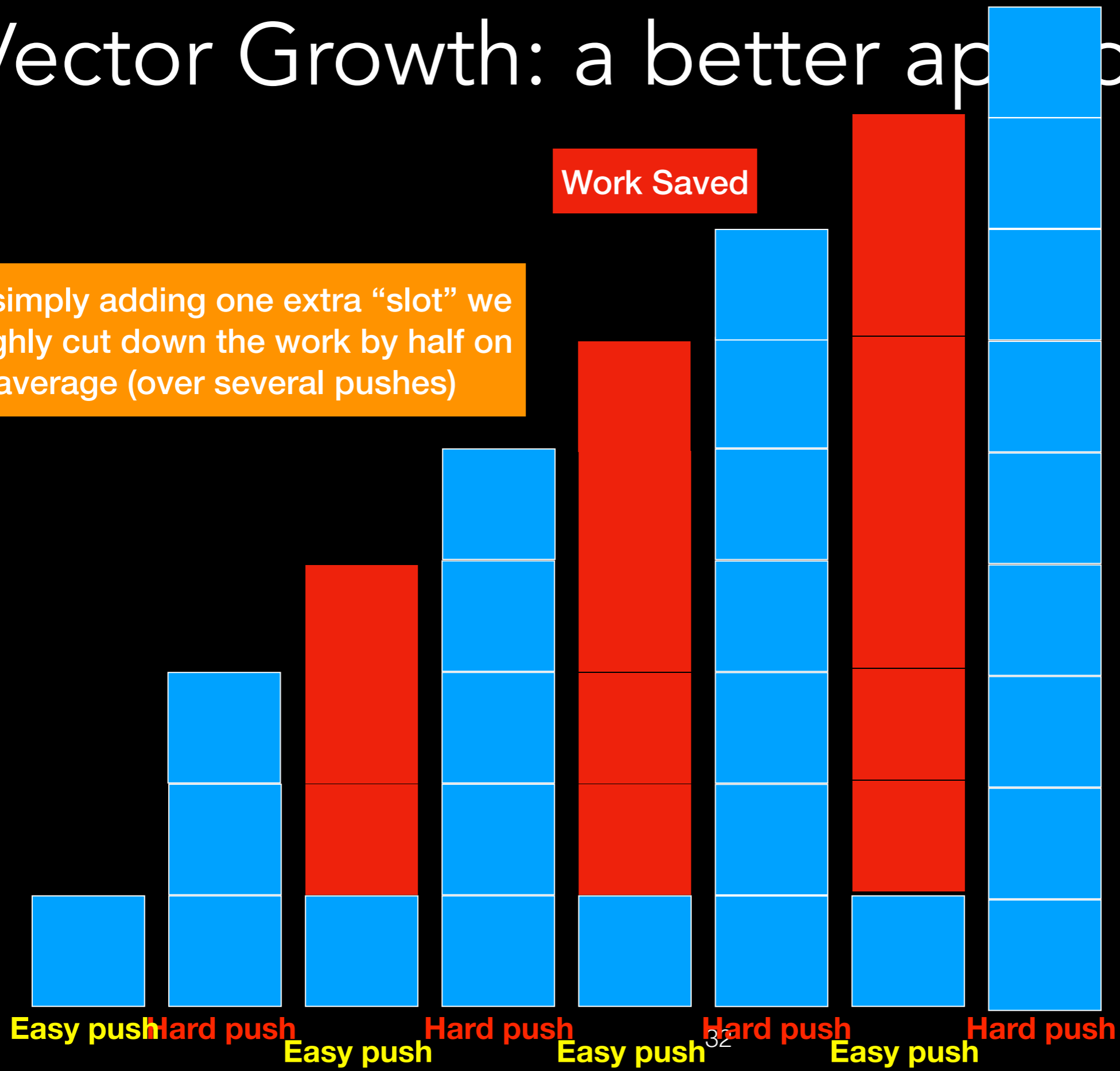
# Vector Growth: a better approach



# Vector Growth: a better approach

Work Saved

By simply adding one extra "slot" we roughly cut down the work by half on average (over several pushes)



Easy push

Hard push

Easy push

Hard push

Easy push

Hard push

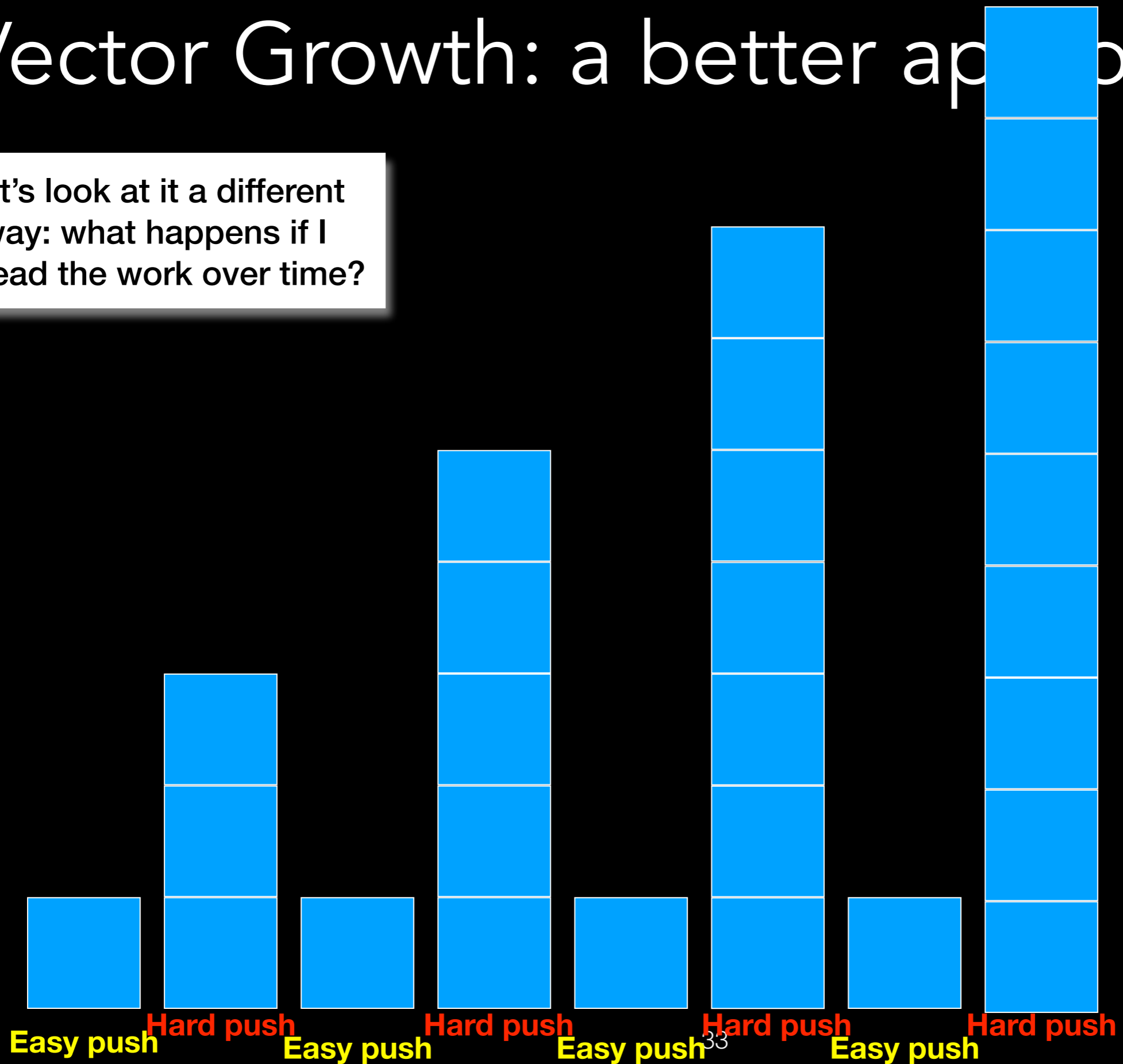
Easy push

Hard push



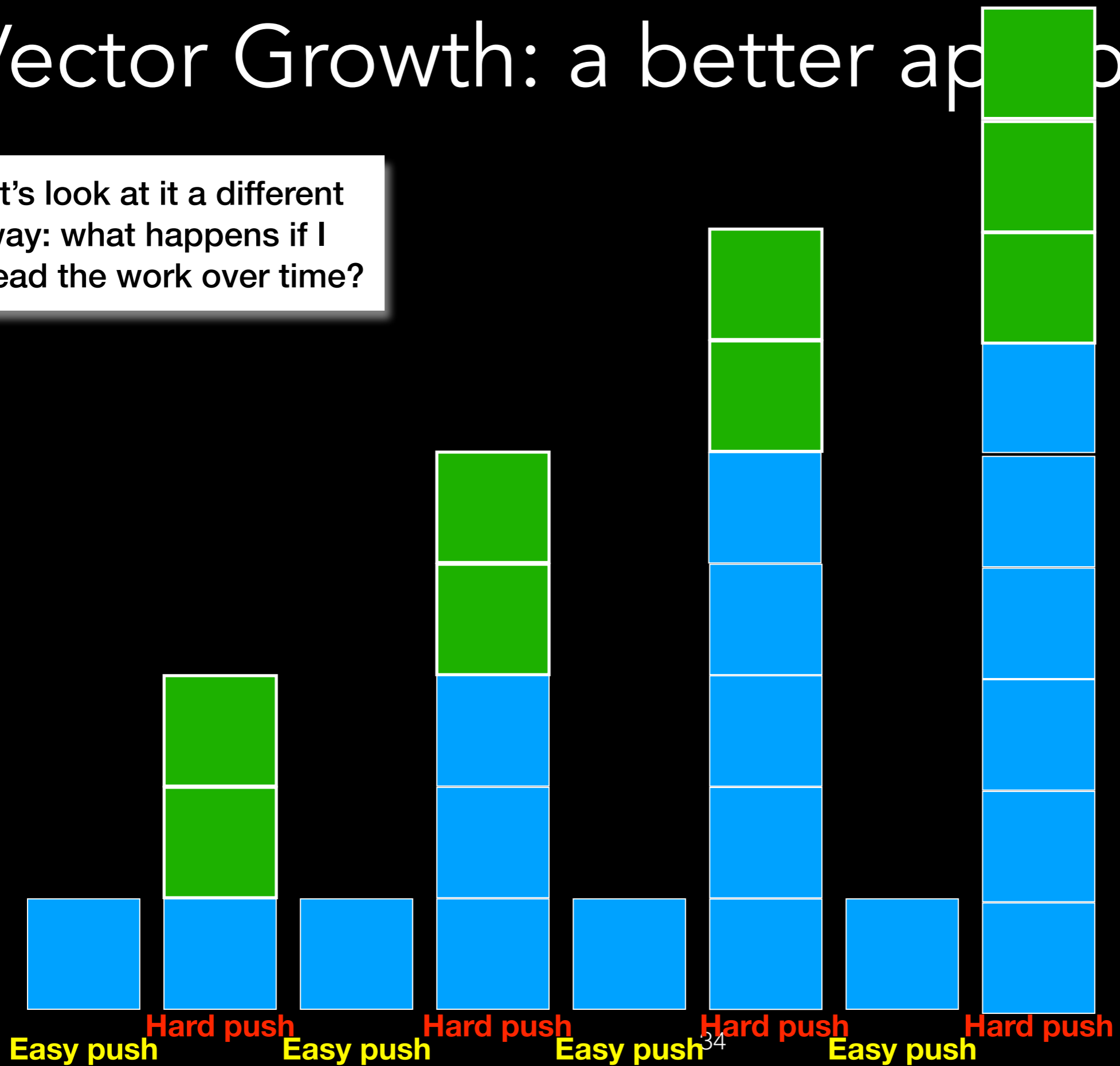
# Vector Growth: a better approach

Let's look at it a different way: what happens if I spread the work over time?



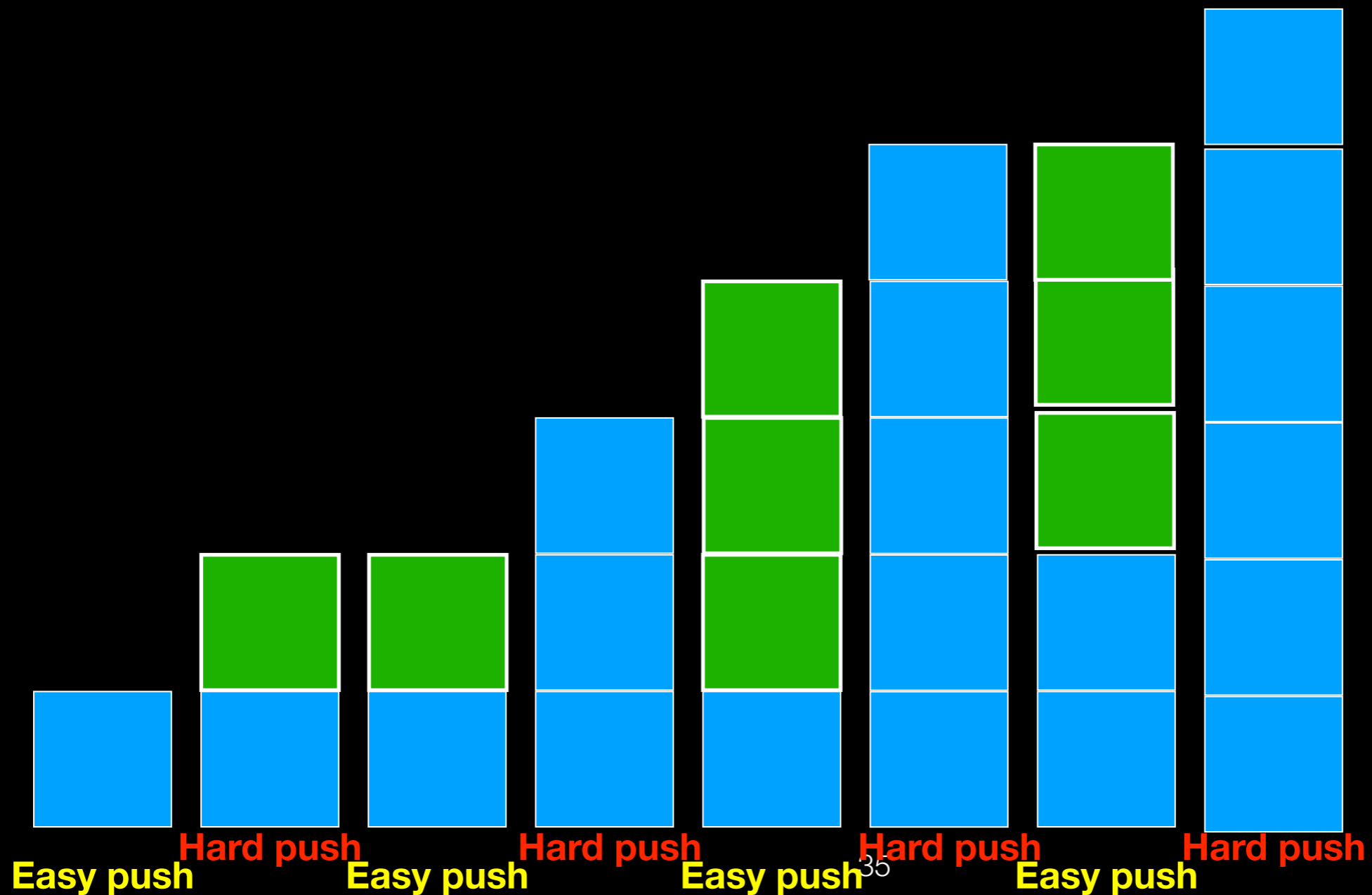
# Vector Growth: a better approach

Let's look at it a different way: what happens if I spread the work over time?



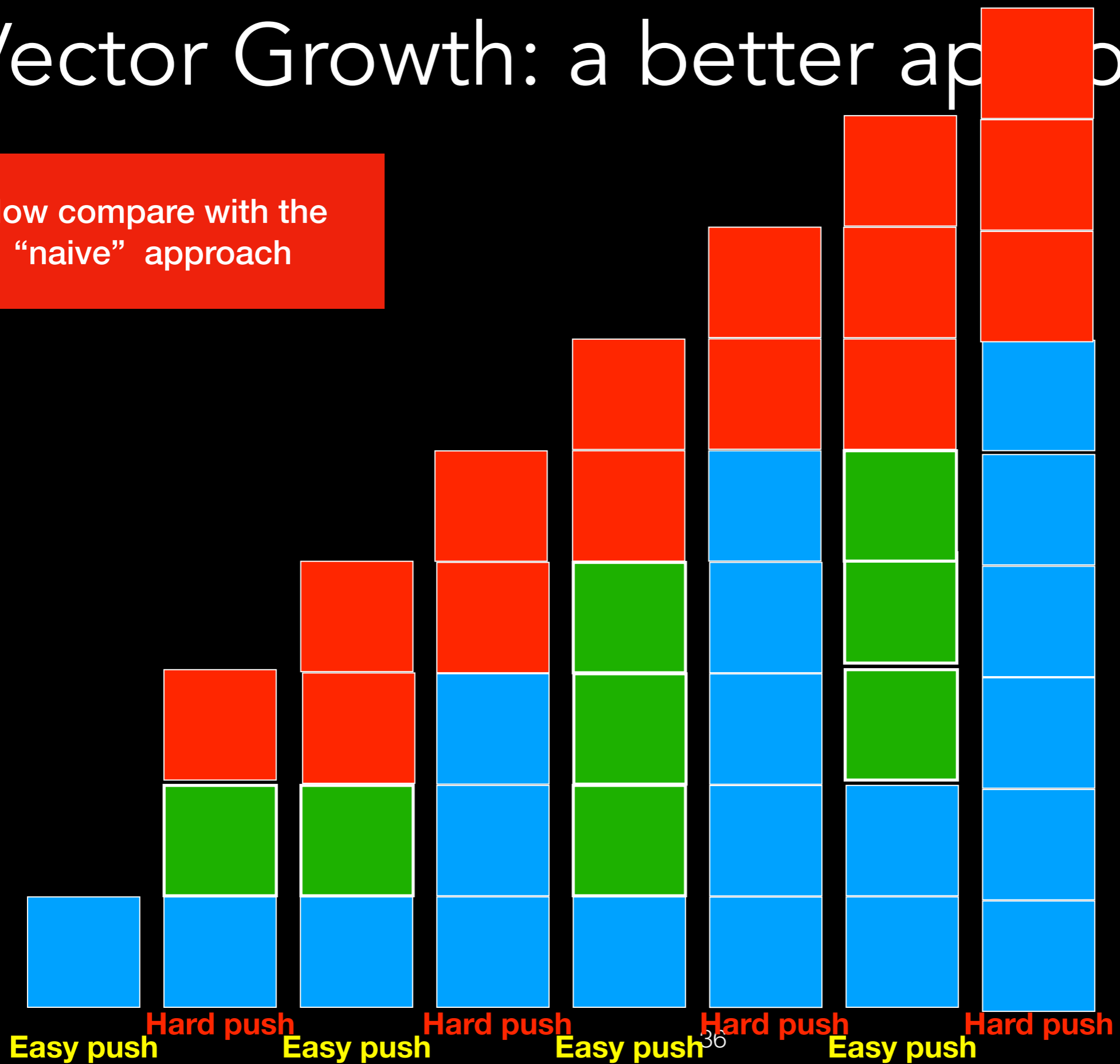
# Vector Growth: a better approach

Let's look at it a different way: what happens if I spread the work over time?



# Vector Growth: a better approach

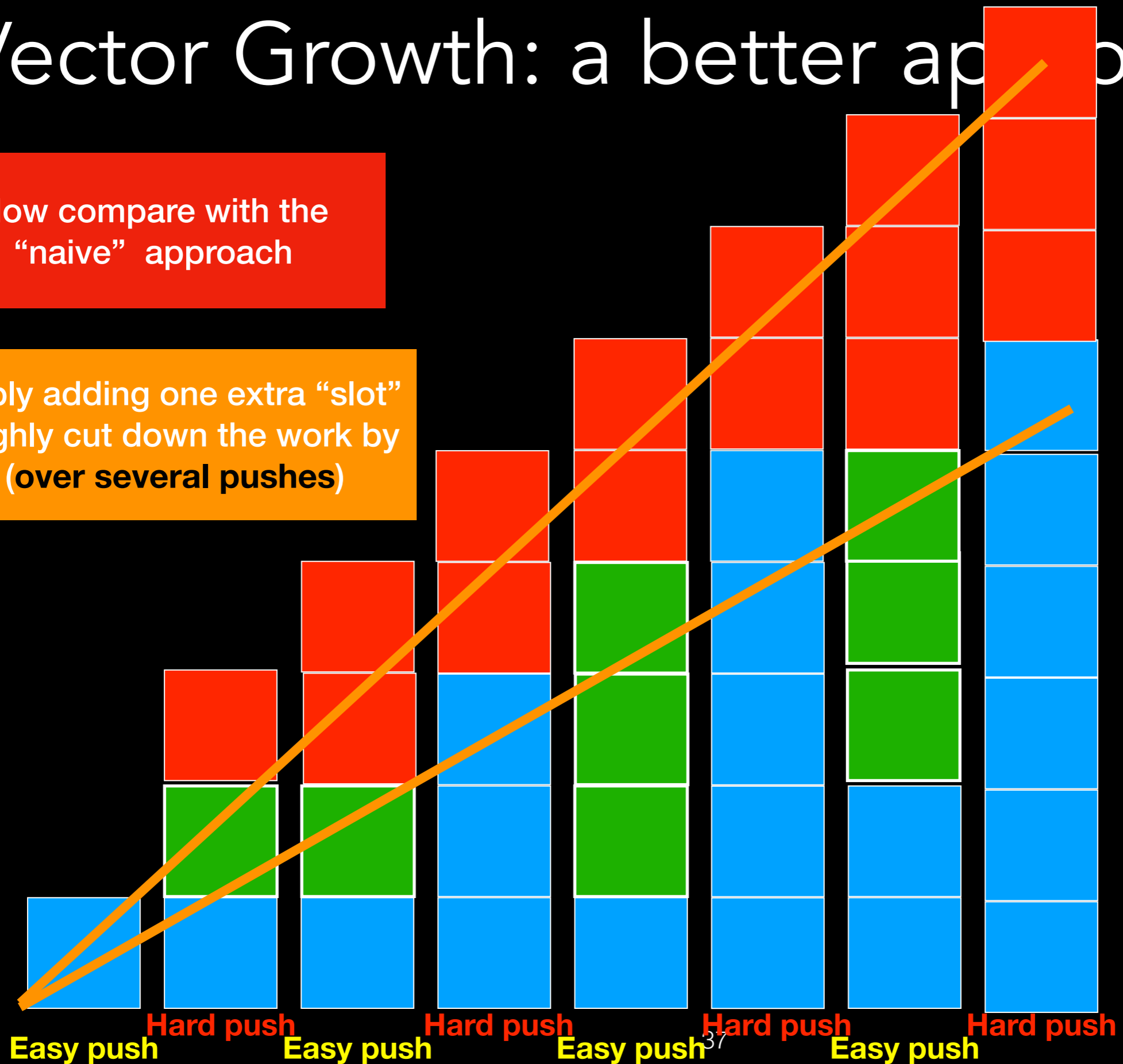
Now compare with the  
“naive” approach



# Vector Growth: a better approach

Now compare with the  
“naive” approach

By simply adding one extra “slot”  
we roughly cut down the work by  
half (over several pushes)



Easy push

Hard push

Easy push

Hard push

Easy push

Hard push

Easy push

Hard push

Can we do better?

# Vector Growth: a much better approach

```
std::vector<T> some_vector;
```

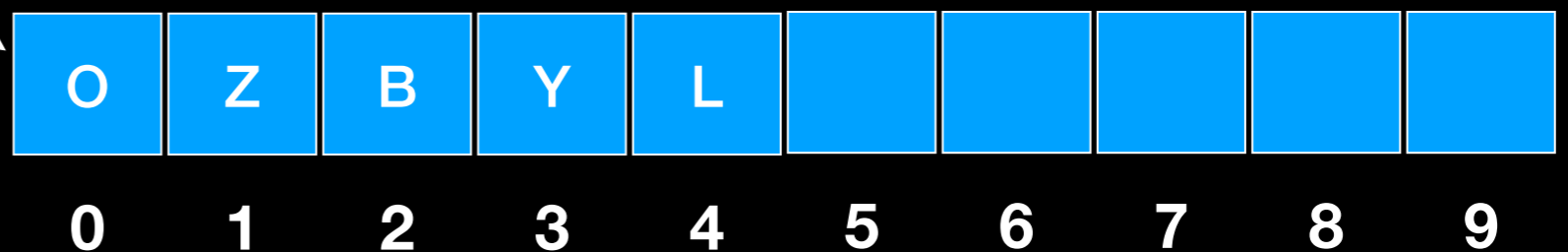
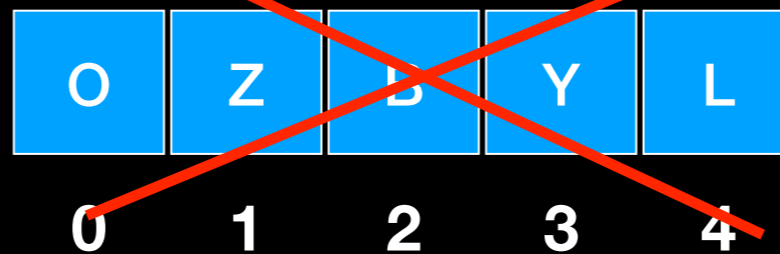
So what is a vector really?



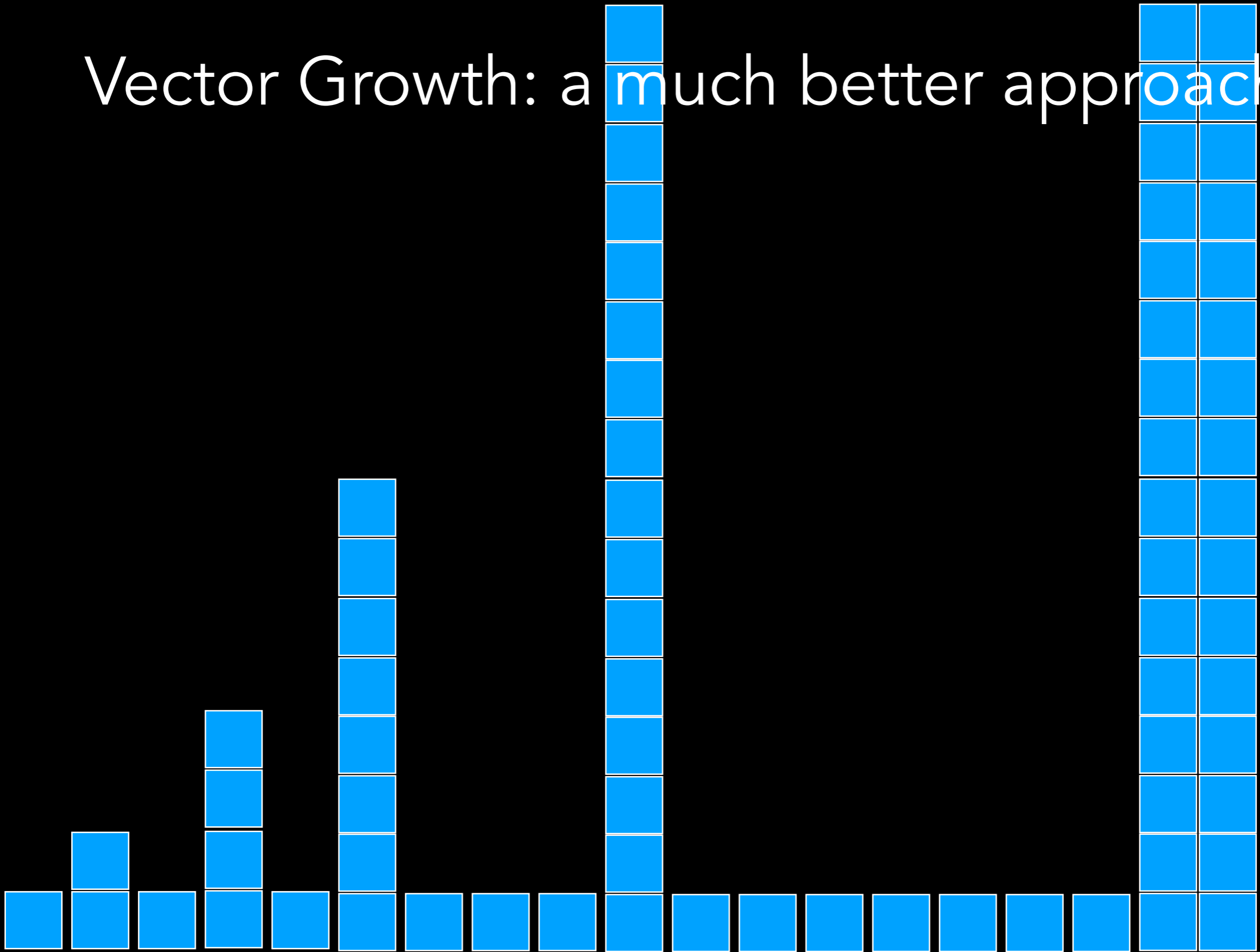
I'll Grow!!!  
I'll double my size!

**Vector (simplified)**

buffer\_ =  
len\_ = 5  
capacity\_ = 10



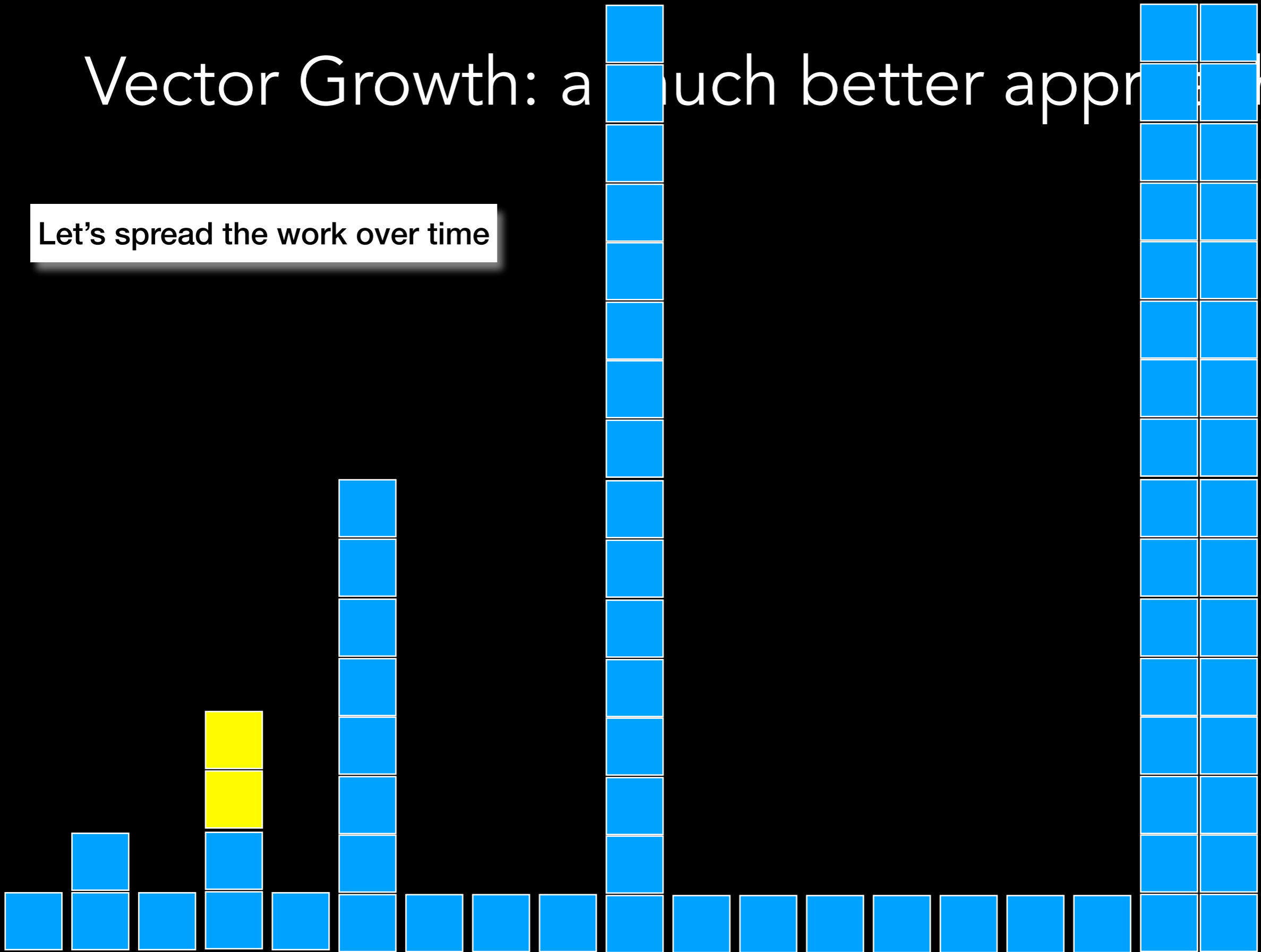
Vector Growth: a much better approach





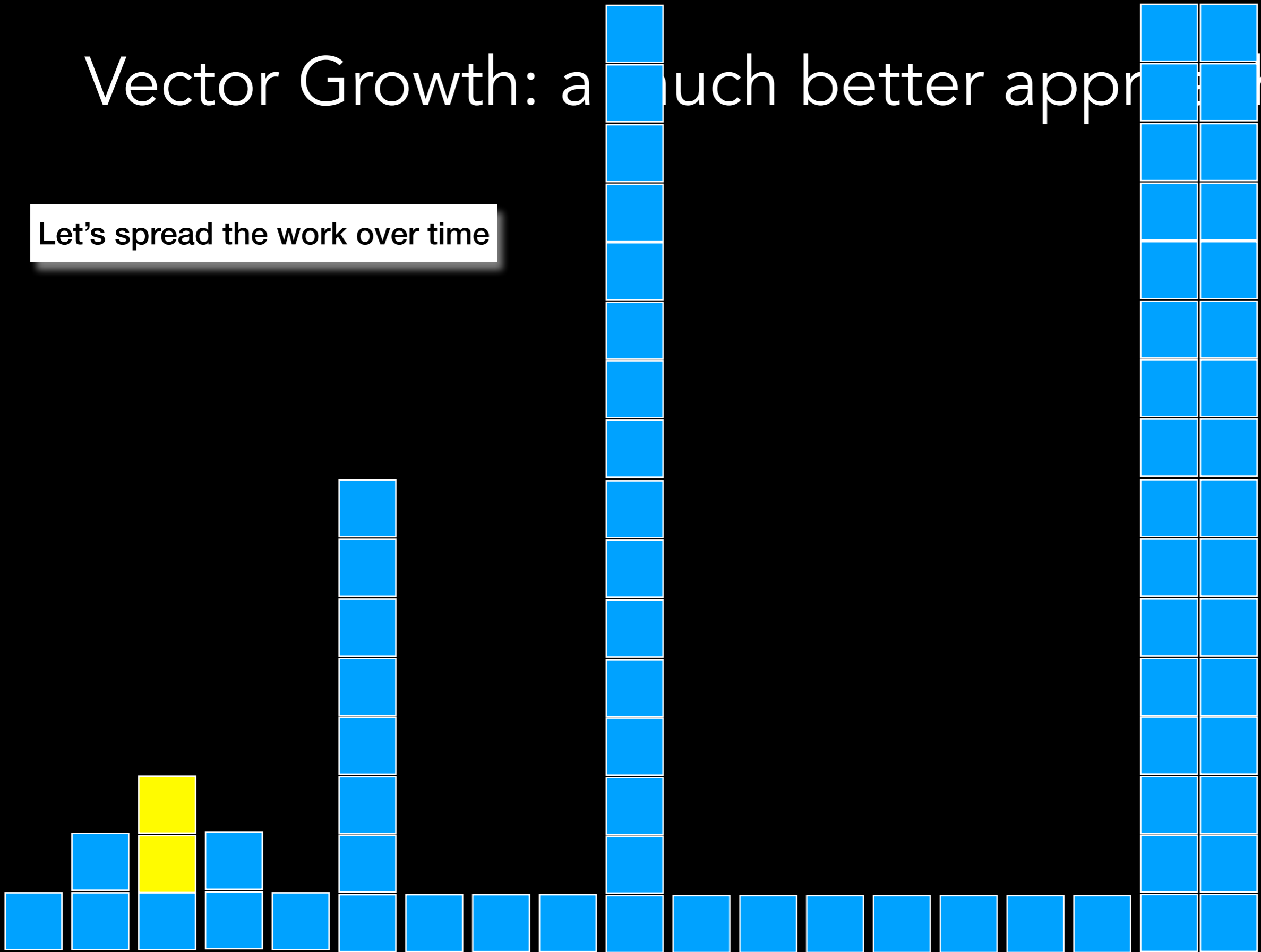
# Vector Growth: a much better approach

Let's spread the work over time



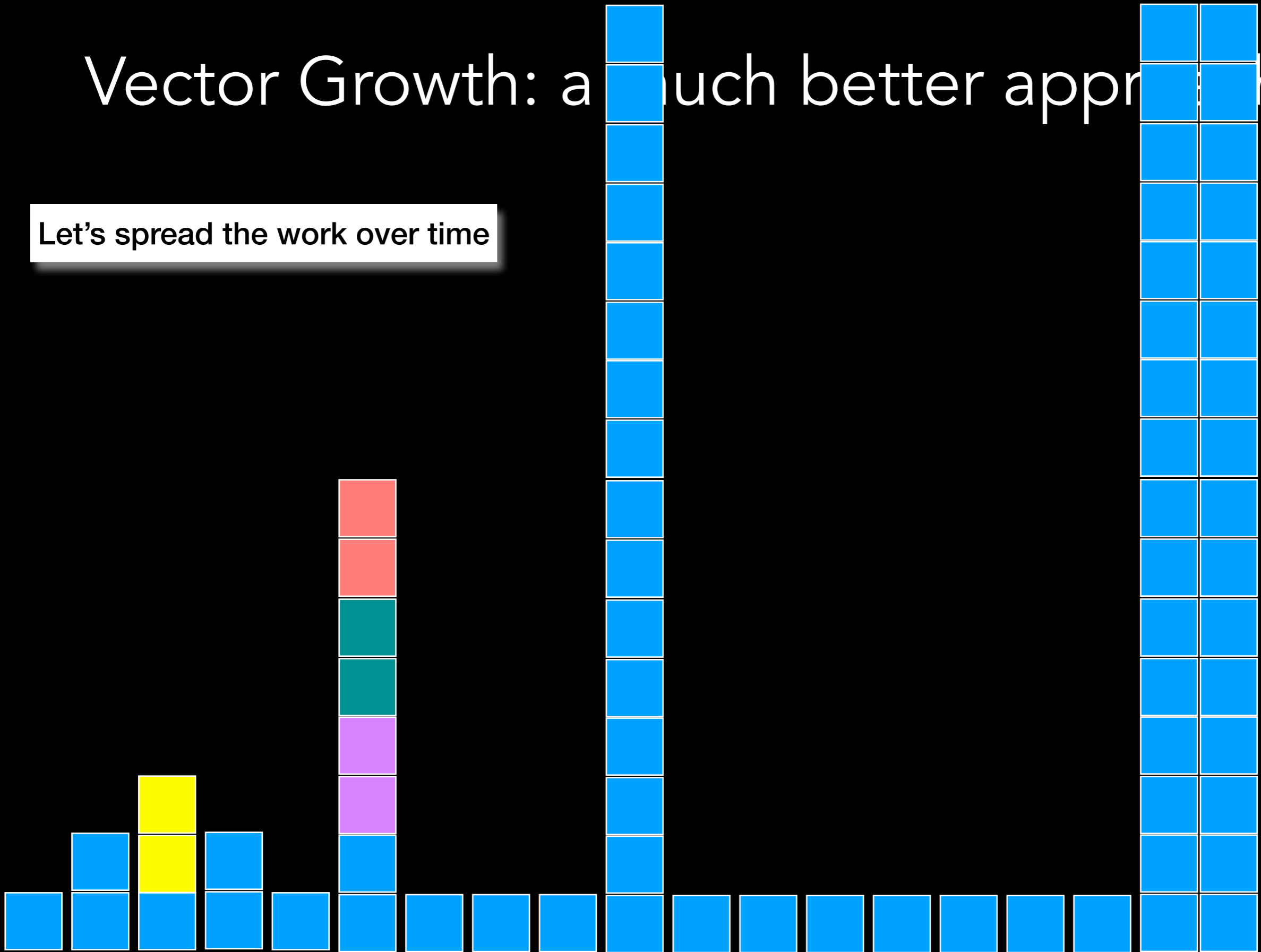
# Vector Growth: a much better approach

Let's spread the work over time



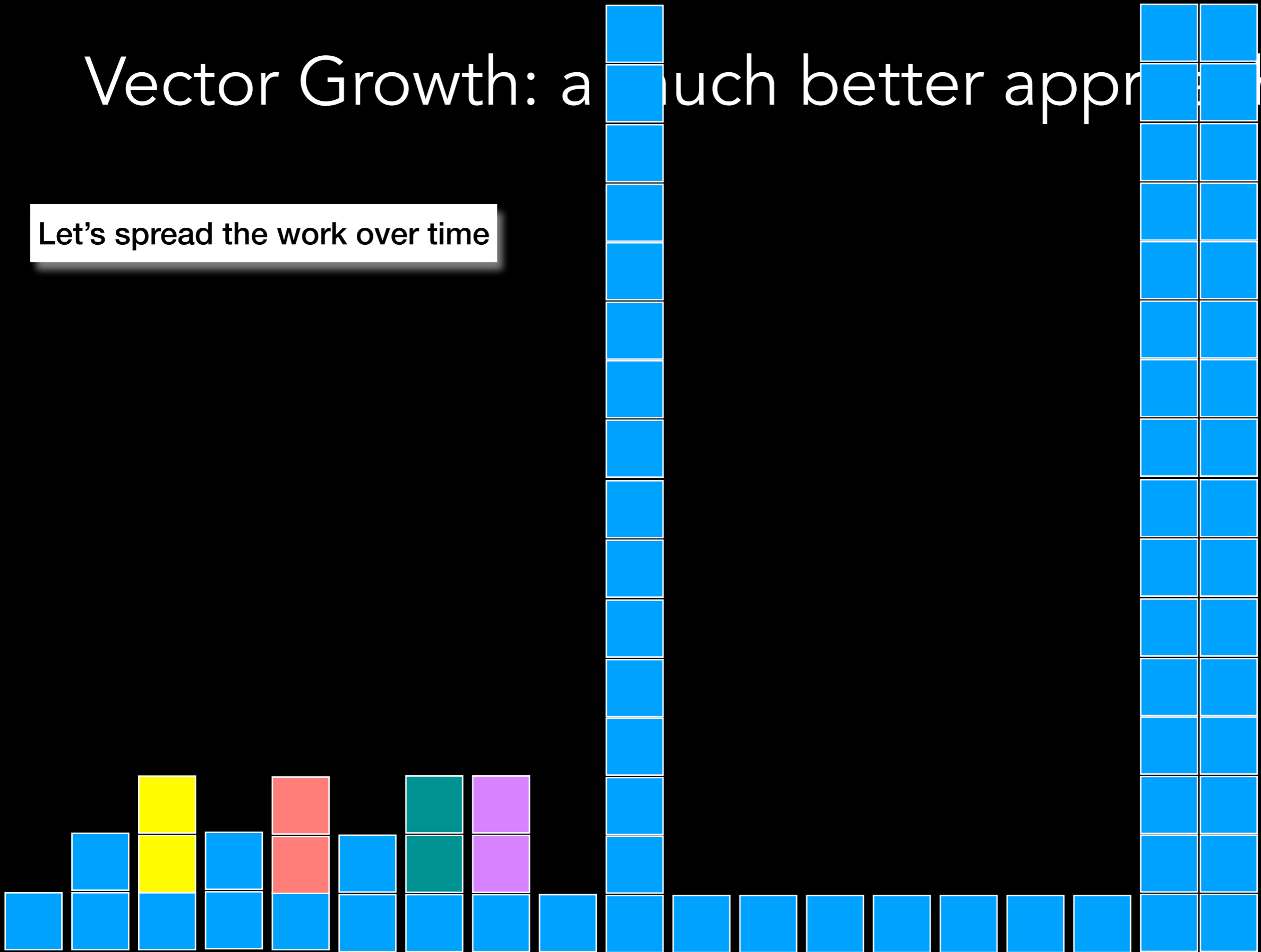
# Vector Growth: a much better approach

Let's spread the work over time



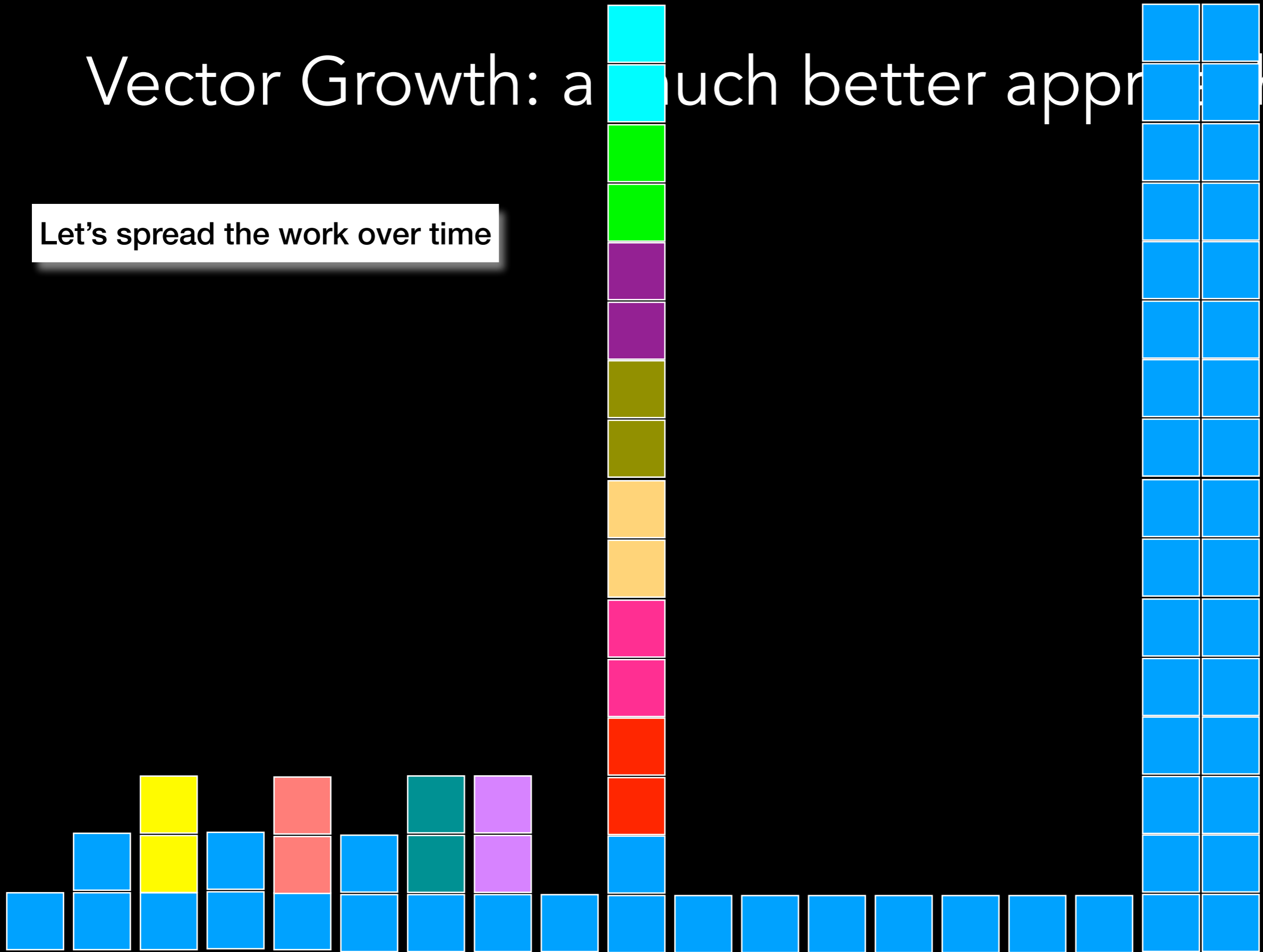
# Vector Growth: a much better approach

Let's spread the work over time



# Vector Growth: a much better approach

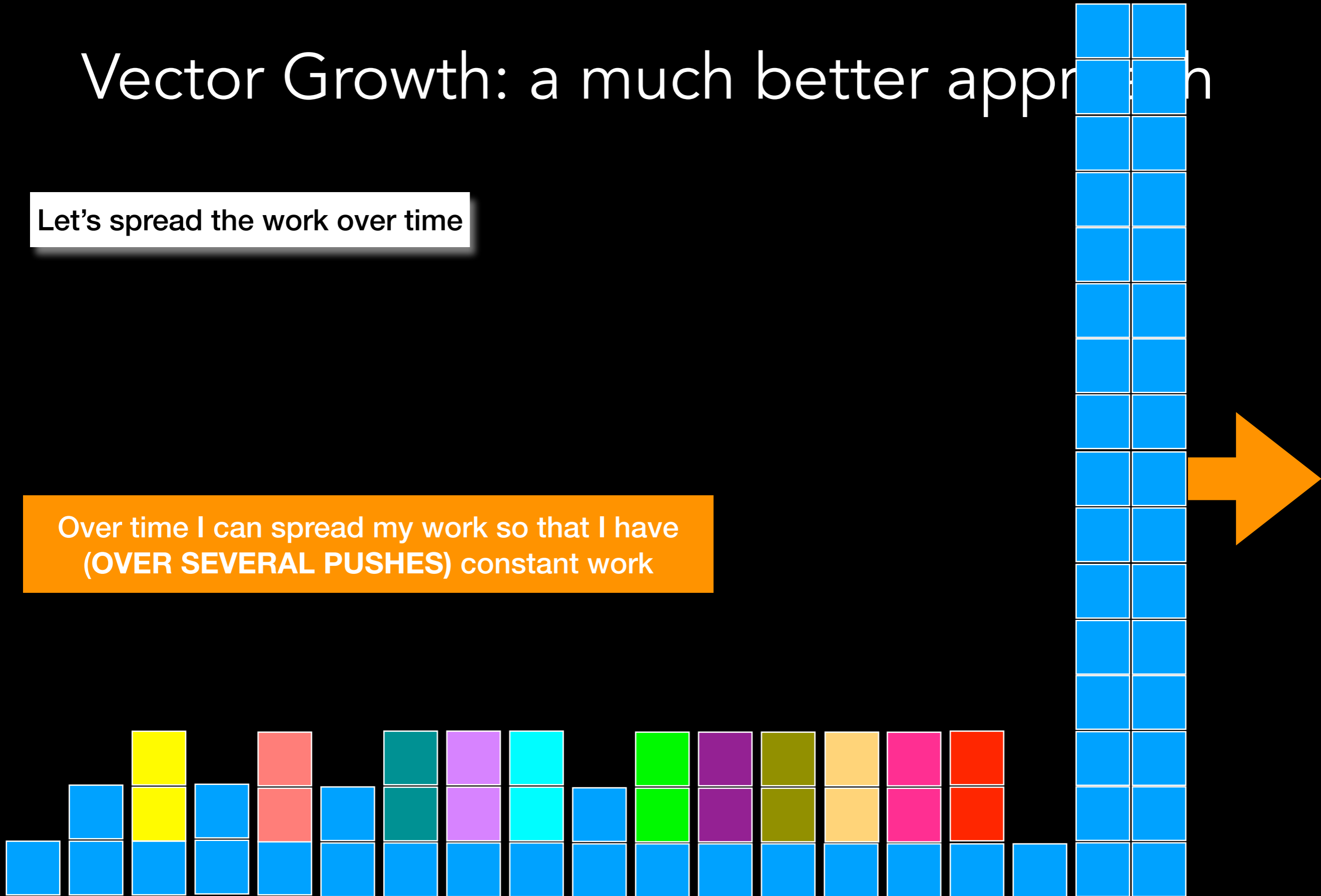
Let's spread the work over time



# Vector Growth: a much better approach

Let's spread the work over time

Over time I can spread my work so that I have  
**(OVER SEVERAL PUSHES)** constant work

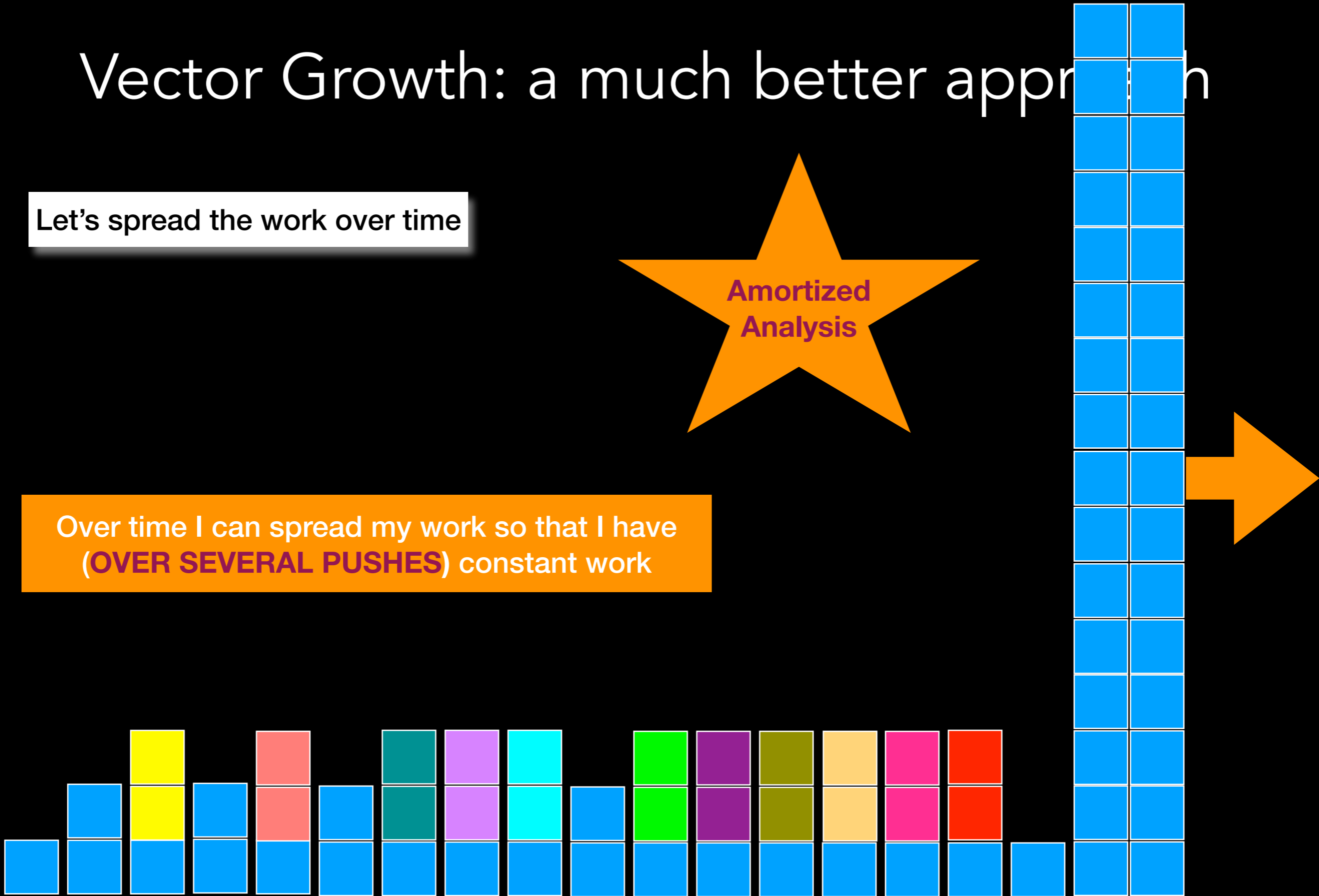


# Vector Growth: a much better approach

Let's spread the work over time

Amortized  
Analysis

Over time I can spread my work so that I have  
**(OVER SEVERAL PUSHES)** constant work



# Vector Growth summarized

If it grows by 1,  $O(n^2)$  over time (**n pushes** - AMORTIZED ANALYSIS)

If it grows by 2, push takes roughly half the "steps" but still  $O(n^2)$  over time (**n pushes** - AMORTIZED ANALYSIS)

If it doubles its size, push takes  $O(1)$  over time (**n pushes** - AMORTIZED ANALYSIS)



# A steadily shrinking Stack

Let's consider this application:

- Push the 524,288<sup>th</sup> ( $2^{19}$ ) element onto Stack which causes it to double its size to 1,048,576 ( $2^{20}$ )
- Reading an input file
  - pop the elements that match
  - manipulate input record accordingly
  - repeat

# A steadily shrinking Stack

Let's consider this application:

- Push the 524,288<sup>th</sup> ( $2^{19}$ ) element onto Stack which causes it to double its size to 1,048,576 ( $2^{20}$ )
- Reading an input file
  - pop the elements that match
  - manipulate input record accordingly
  - repeat

How much I pop will depend on input

# A steadily shrinking Stack

Let's consider this application:

- Push the 524,288<sup>th</sup> ( $2^{19}$ ) element onto Stack which causes it to double it's size to 1,048,576 ( $2^{20}$ )
- Reading an input file
  - pop the elements that match
  - manipulate input record accordingly
  - repeat

Assume a few matches at each iteration -> mostly empty stack but it will be around for a long time!



I will not shrink!

Useless  
memory  
waste

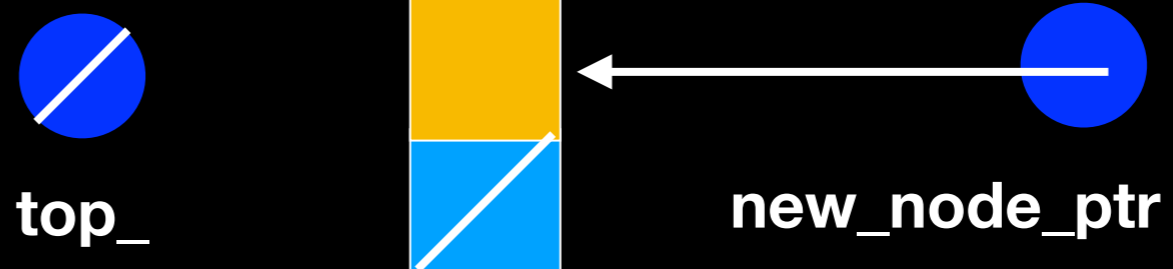
# Linked Chain



top\_

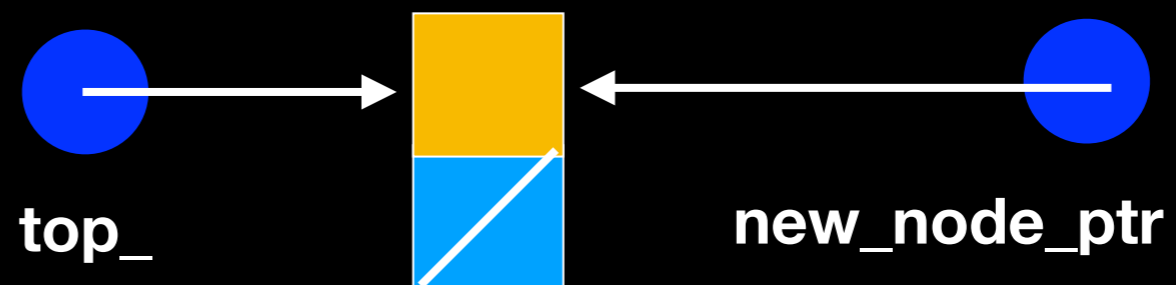
# Linked Chain

push



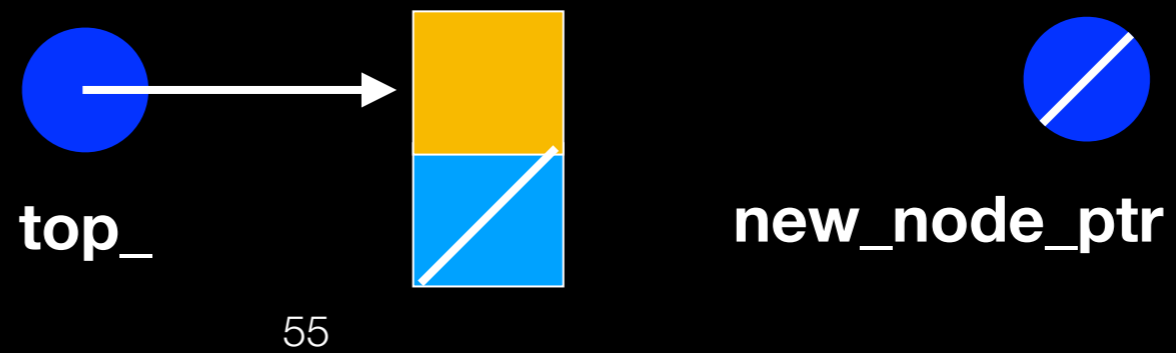
# Linked Chain

push

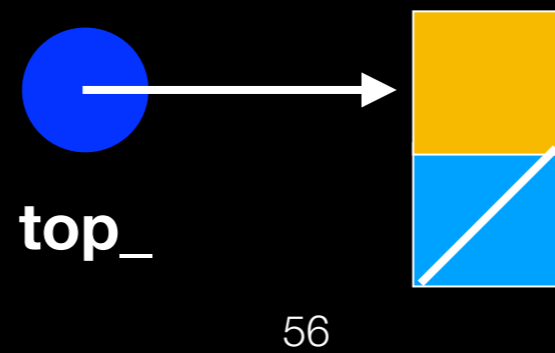


# Linked Chain

push



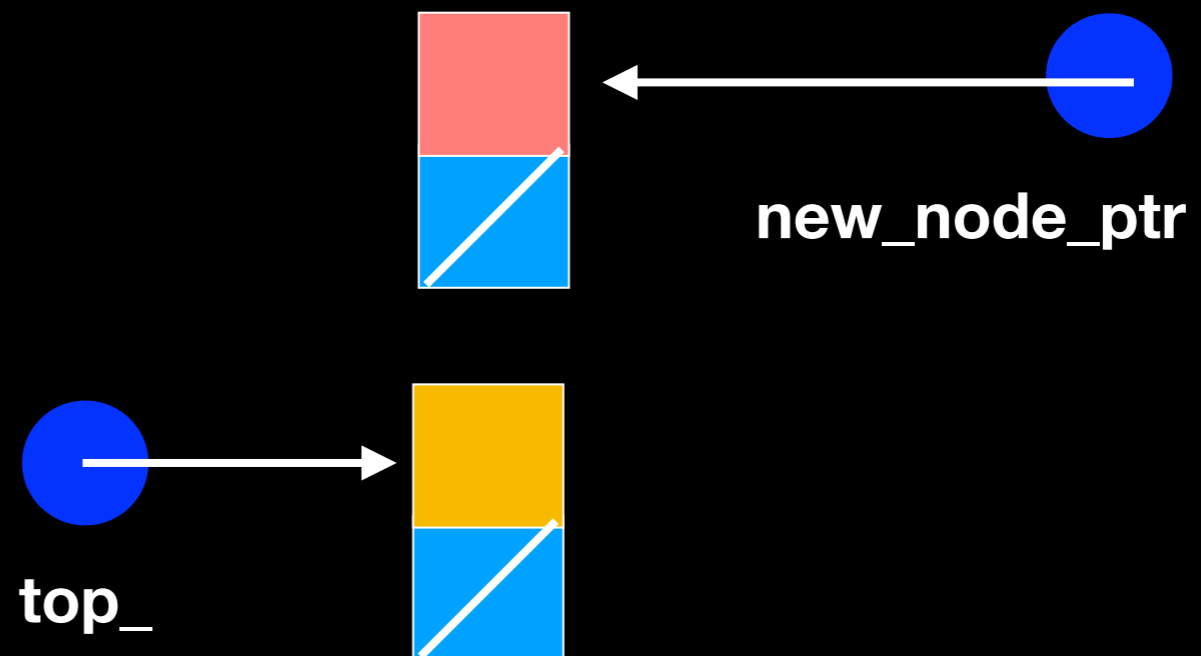
# Linked Chain





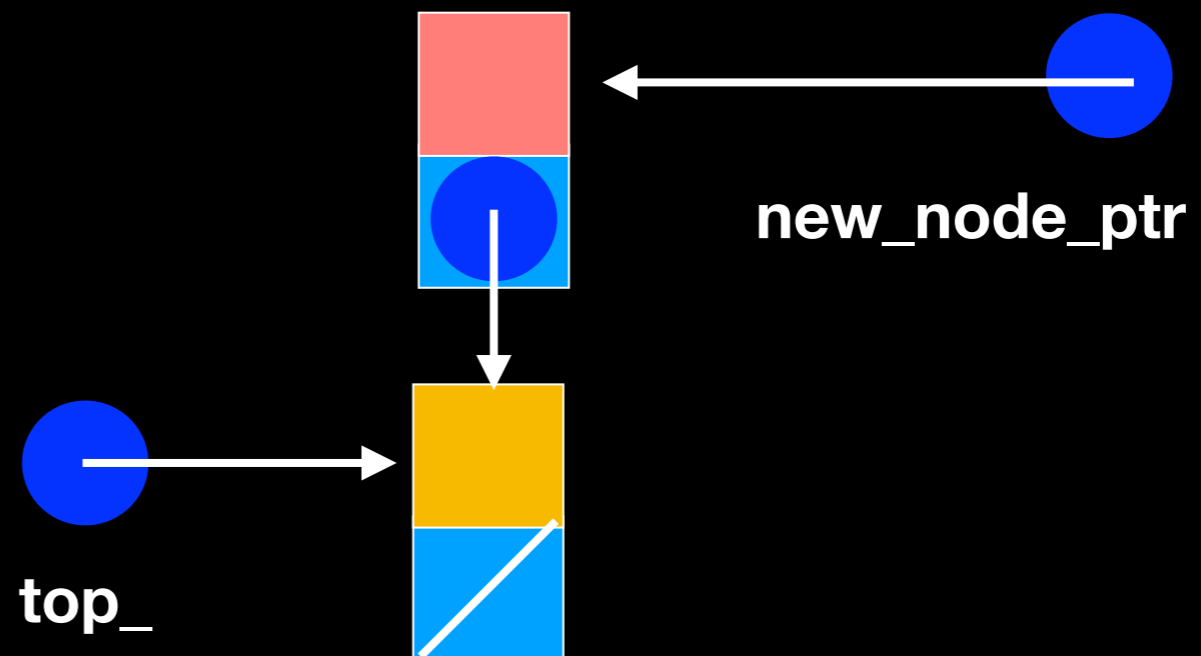
# Linked Chain

push



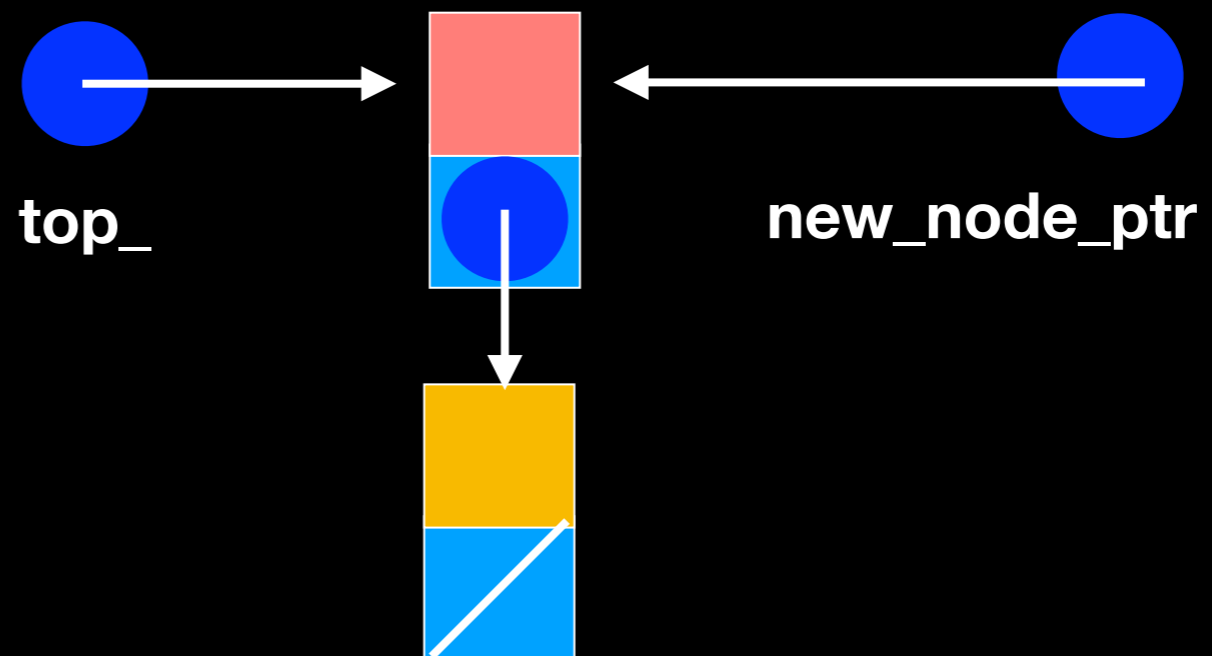
# Linked Chain

push



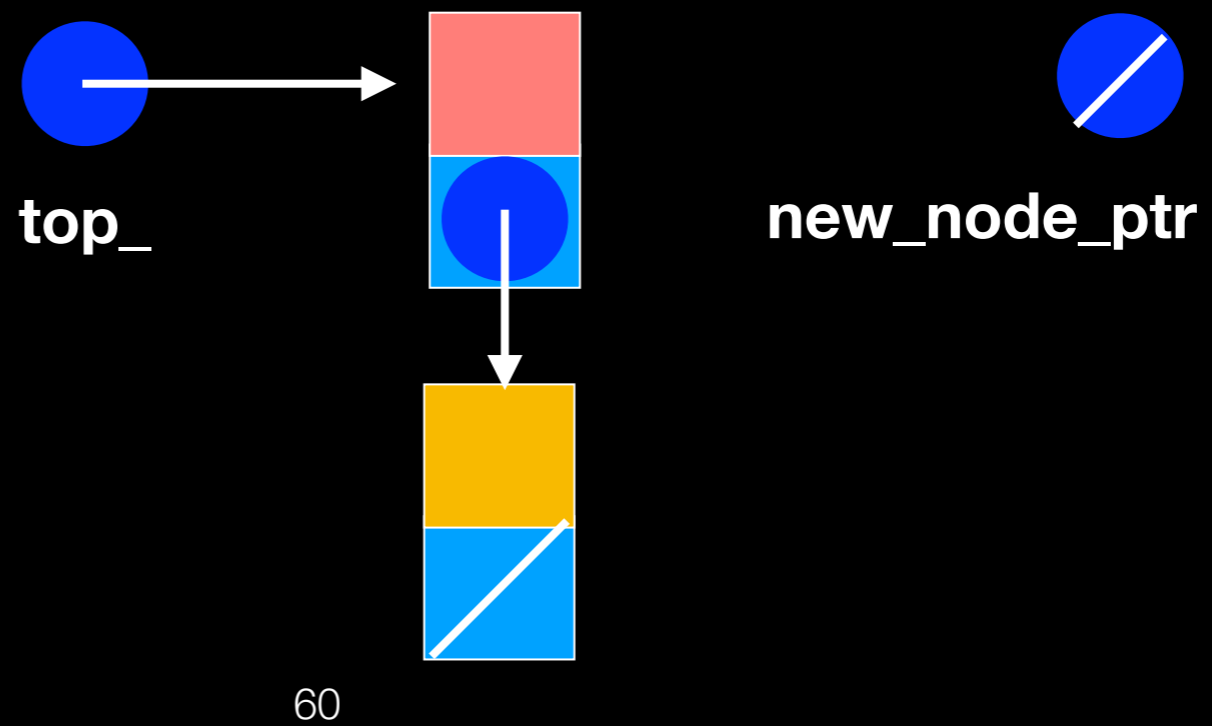
# Linked Chain

push

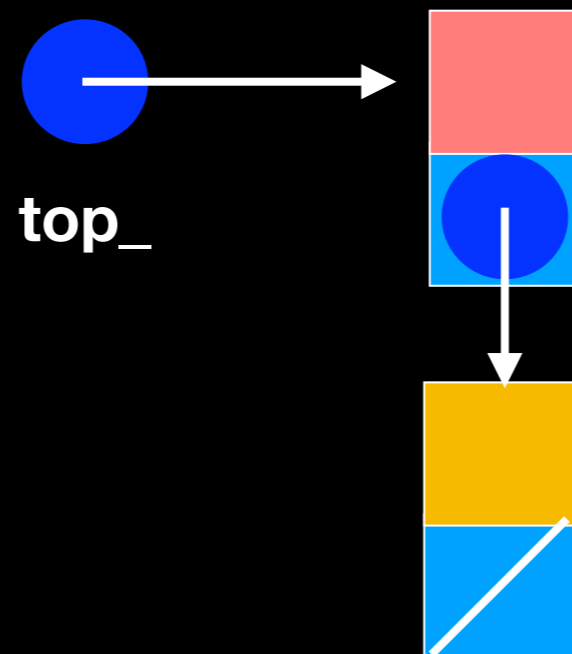


# Linked Chain

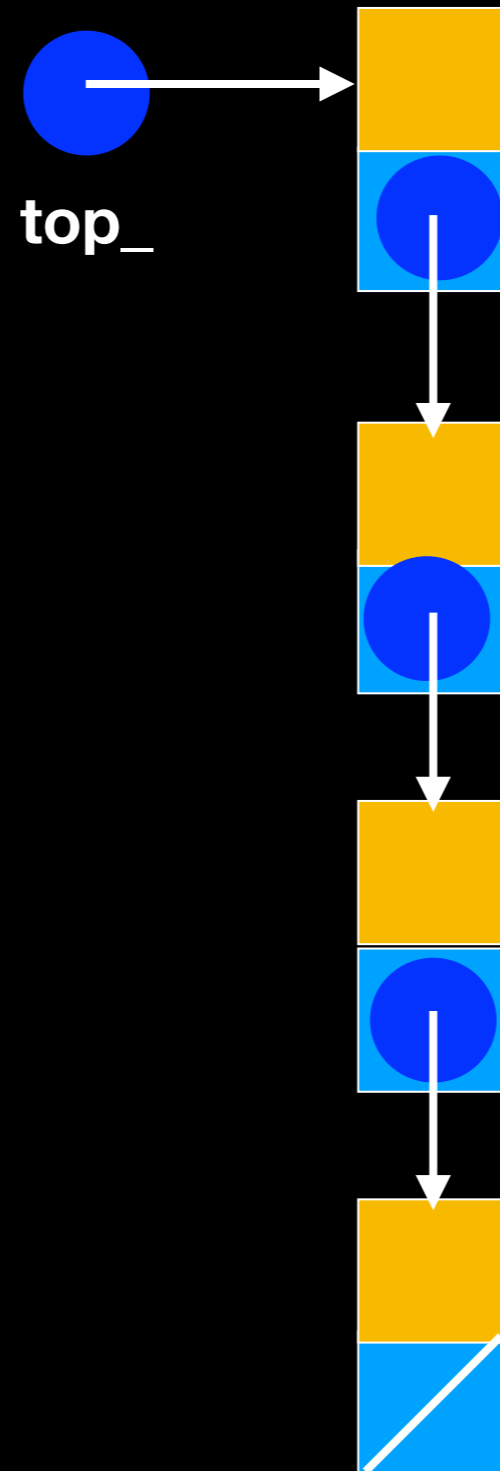
push



# Linked Chain

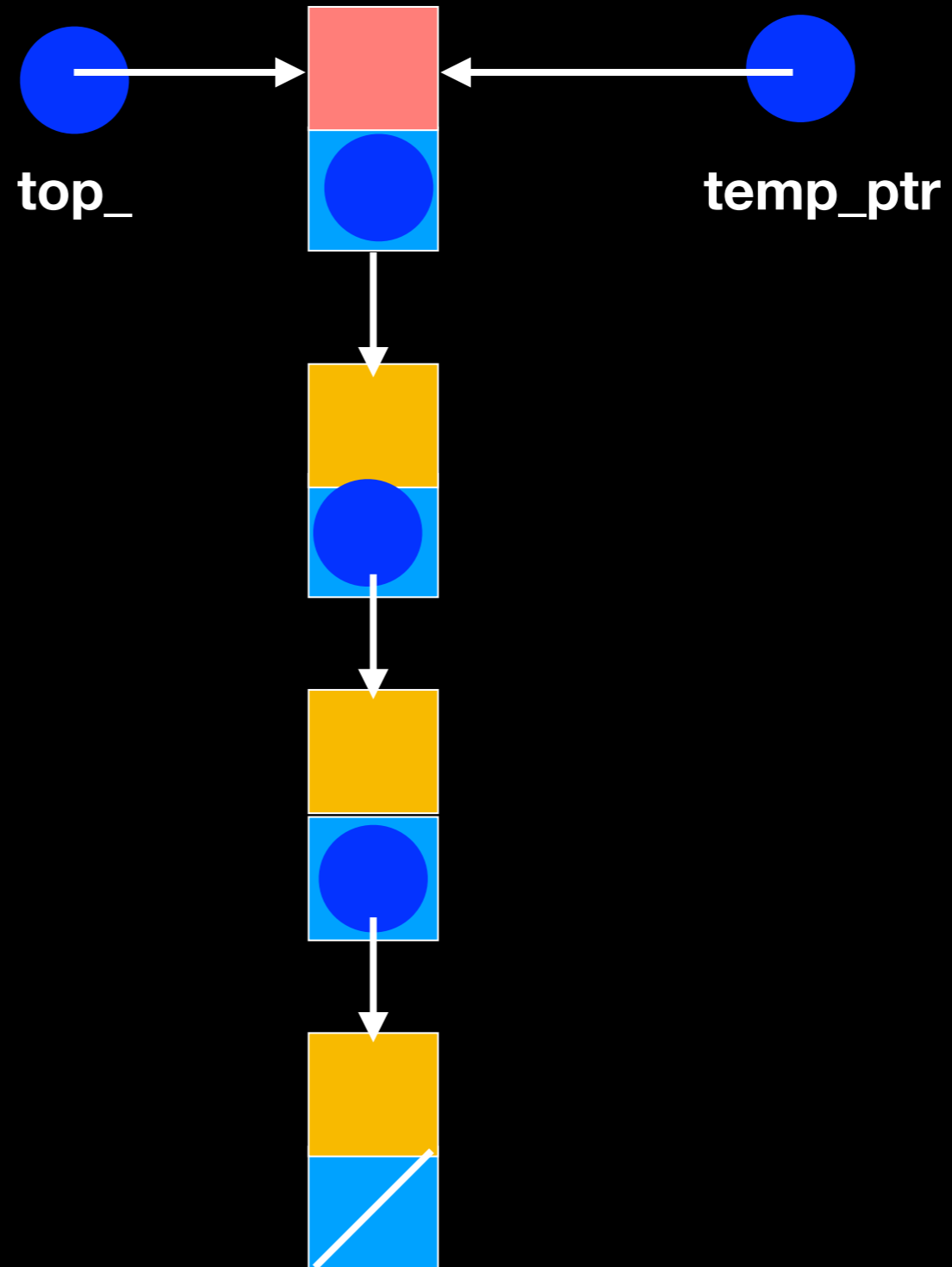


# Linked Chain



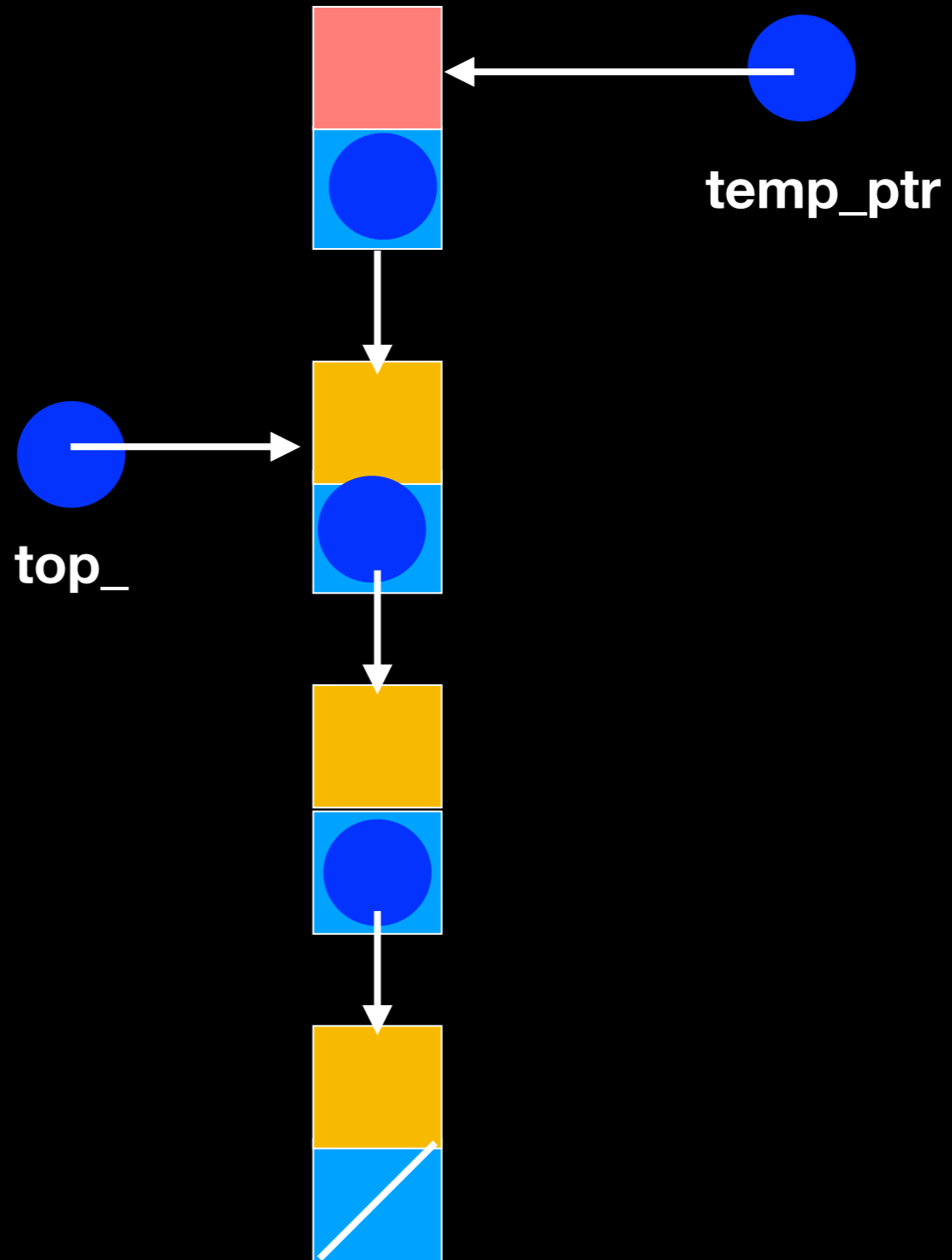
# Linked Chain

pop



# Linked Chain

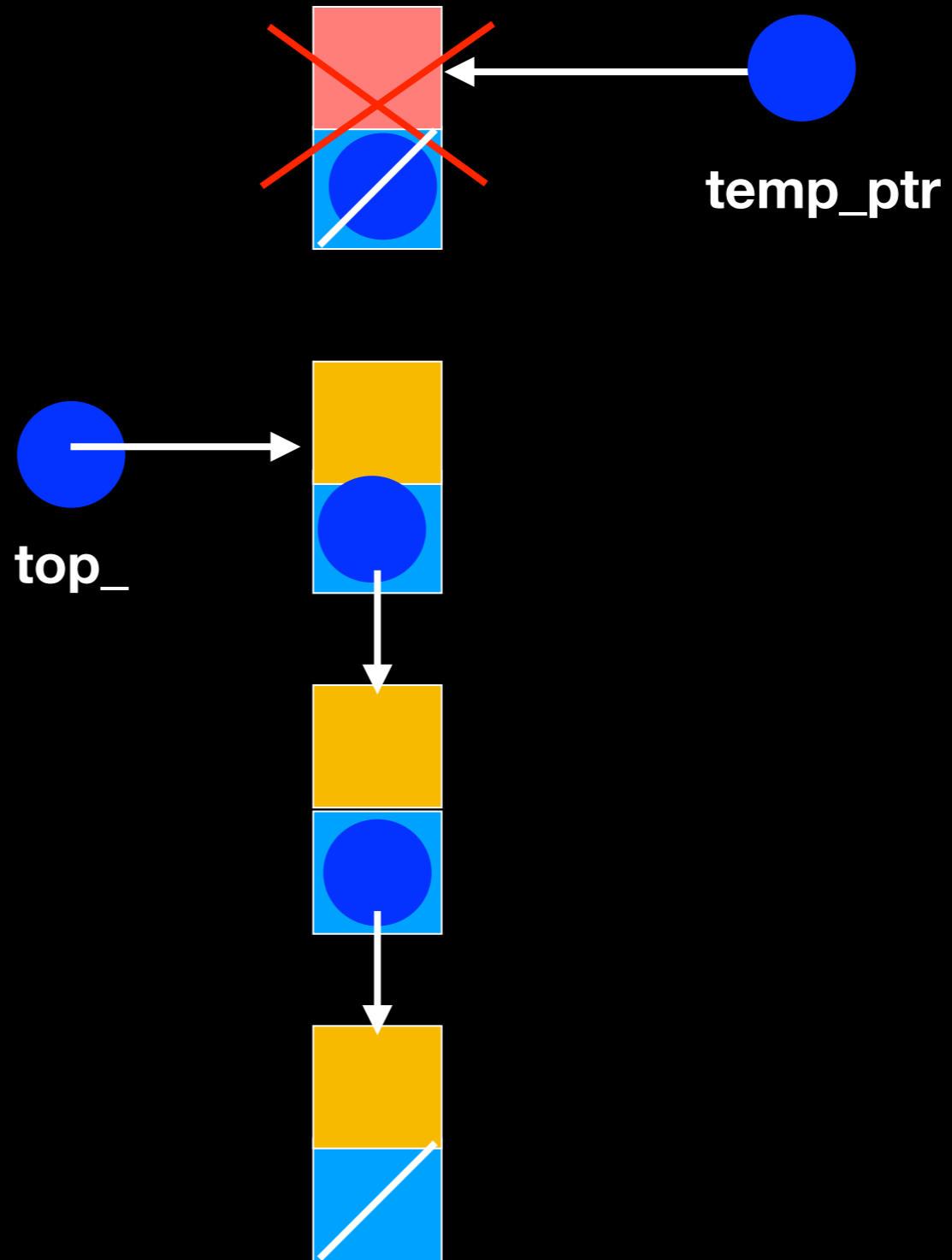
pop





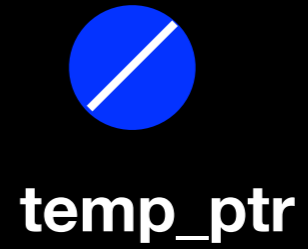
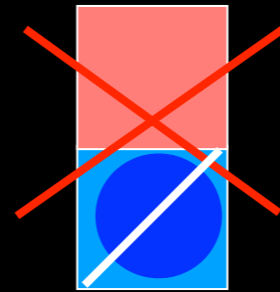
# Linked Chain

pop

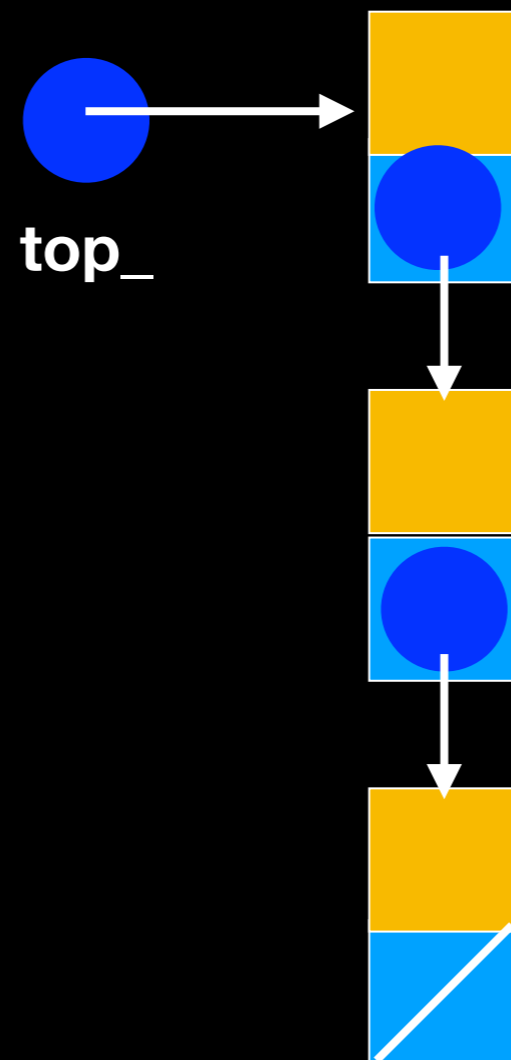


# Linked Chain

pop



# Linked Chain



# Linked-Chain Analysis

1 assignment + 1 increment/decrement =  $O(1)$

*size* :  $O(1)$

*isEmpty*:  $O(1)$

*push*:  $O(1)$

*pop* :  $O(1)$

*top* :  $O(1)$

GREAT!!!! And there is no "Except" case here, every operation is  $O(1)$ !

# To summarize

**Array:**  $O(1)$  for push and pop, but **size is bounded**

**Vector:** size is unbounded but

-Some push operations take  $O(1)$ , others take

$O(n) \rightarrow O(1)$  over time (AMORTIZED ANALYSIS)

**Linked-Chain:**  $O(1)$  for push and pop and **size is unbounded**

# Implement Stack ADT

```
#ifndef STACK_H_
#define STACK_H_

template<typename ItemType>
class Stack
{
public:
    Stack();
    void push(const ItemType& new_entry); //adds an element to top of stack
    void pop(); // removes element from top of stack
    ItemType top() const; // returns a copy of element at top of stack
    int size() const; // returns the number of elements in the stack
    bool isEmpty() const; //returns true if no elements on stack, else false

private:
    //implementation details here

}; //end Stack

#include "Stack.cpp"
#endif // STACK_H_`
```

What should we add here to implement it as a linked chain?

# Implement Stack ADT

```
#ifndef STACK_H_  
#define STACK_H_
```

```
template<class T>  
class Stack  
{
```

```
public:
```

```
Stack();
```

```
~Stack(); // destructor
```

```
Stack(const Stack<T>& a_stack); //copy constructor
```

```
void push(const T& newEntry); // adds an element to top of stack
```

```
void pop(); // removes element from top of stack
```

```
T top() const; // returns a copy of element at top of stack
```

```
int size() const; // returns the number of elements in the stack
```

```
bool isEmpty() const; //returns true if no elements on stack else false
```

```
private:
```

```
Node<T>* top_; // Pointer to top of stack
```

```
int item_count; // number of items currently on the stack
```

```
}; //end Stack
```

```
#include "Stack.cpp"
```

```
#endif // STACK_H_
```